



IPv6 Security

SEC-2003

Cisco Networkers
2006

Session Objectives

- **This session presents IPv6 security as contrasted with IPv4 from a threat and mitigation perspectives**
- **Advanced IPv6 security topics like transition options and dual stack IPv6/IPv4 environments**
- **This session requires a basic knowledge of the IPv6 and IPsec protocols as well as IPv4 security best practices**

Agenda

- **Types of Threats**
- **IPv6 and IPv4 Threat Comparisons**
- **IPv6 Security Best Common Practice**
- **Specific Issues for IPv6**
 - Tunnels and Mobile IPv6**
- **Enforcing a Security Policy in IPv6**
 - ACL and Firewalls**
- **Enterprise Secure Deployment**

Types of Threats



Types of Threats

- **Reconnaissance**—Provide the adversary with information
- **Unauthorized access**—Exploit
- **Header manipulation and fragmentation**—Evade or overwhelm
- **Layer 3–Layer 4 spoofing**— Mask the intent or origin of the traffic
- **ARP and DHCP attacks**—Subvert the host initialization process
- **Broadcast amplification attacks (smurf)**—Amplify the effect of a flood
- **Routing attacks**—Disrupt or redirect traffic flows

Types of Threats (Cont.)

- **Viruses and worms**— Propagation of the malicious payload
- **Sniffing**—Capturing data
- **Application layer attacks**— Attacks executed at Layer 7
- **Rogue devices**—Unauthorized devices connected to a network
- **Man-in-the-middle attacks**— Attacks which involve interposing an adversary between two communicating parties
- **Flooding**—Consume enough resources to delay processing of valid traffic

Threat Comparison



Reconnaissance in IPv4

In IPv4, Reconnaissance Is Relatively Easy

1. DNS/IANA crawling (whois) to determine ranges
2. Ping sweeps and port scans
3. Application vulnerability scans

```
[tick:/var] scott# nmap -sP 10.1.1.0/24
Starting nmap V. 3.00 ( www.insecure.org/nmap/ )
Host (10.1.1.0) seems to be a subnet broadcast ...
Host (10.1.1.1) appears to be up.
Host (10.1.1.12) appears to be up.
Host (10.1.1.22) appears to be up.
Host (10.1.1.23) appears to be up.
Host (10.1.1.101) appears to be up.
Host (10.1.1.255) seems to be a subnet broadcast ...
Nmap run completed -- 256 IP addresses (7 hosts up)
scanned in 4 seconds
```


Reconnaissance in IPv6

Subnet Size Difference

- **Default subnets in IPv6 have 2^{64} addresses**
=> scanning every address: centuries vs. seconds
- **NMAP doesn't even support for ping sweeps on IPv6 networks (you'll have retired by the time it finishes, even at one million packets per second)**

Reconnaissance in IPv6

IP6 Scanning Methods Are Likely to Change

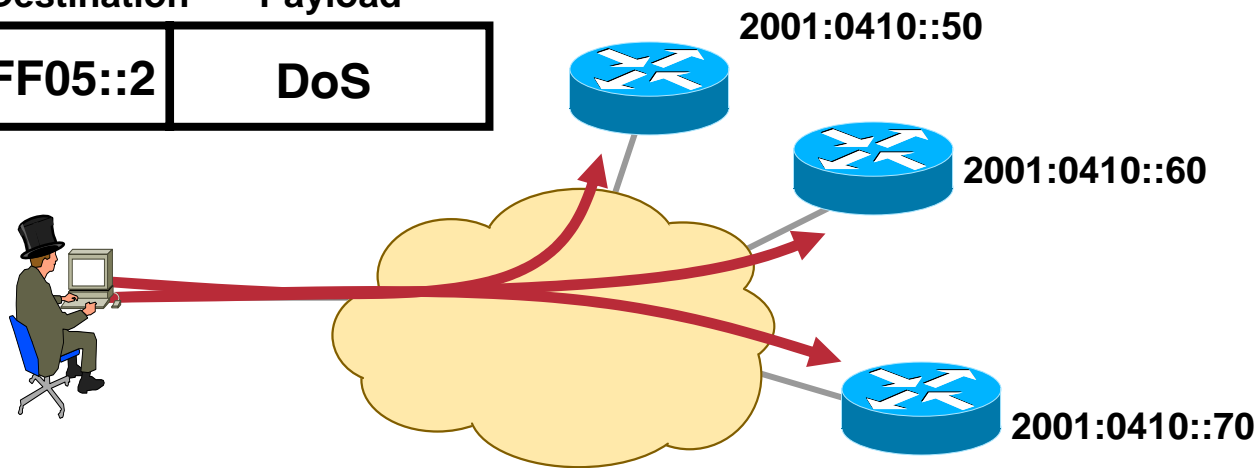
- **Public servers will still need to be DNS reachable**
- **Dynamic DNS adoption causing DNS servers to be rich sources of addresses to scan**
- **Administrators may adopt easy to remember addresses (::10,::20,::F00D, or simply IPv4 last octet)**
- **By compromising hosts in a network, an attacker can learn new addresses to scan**

Reconnaissance in IPv6

New Multicast Addresses

- For example, all routers (FF05::2) and all DHCP servers (FF05::1:3)
- These addresses must be filtered at the border in order to make them unreachable from the outside

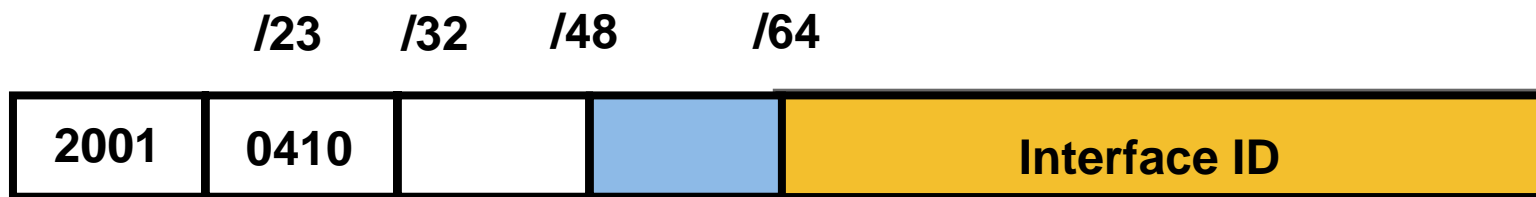
Source	Destination	Payload
Attacker	FF05::2	DoS



Reconnaissance IPv6 Best Practices

- **Implement privacy extensions carefully—**
(next slide)
- **Filter internal-use IPv6 addresses at organization border routers—**prevent addresses like the all nodes multicast address from becoming conduits for attack
- **Filter unneeded services at the firewall—**just like in IPv4
- **Selectively filter ICMP—**more on this later
- **Maintain host and application security—**just like in IPv4

IPv6 Privacy Extensions (RFC 3041)



- **Temporary addresses for IPv6 host client application, e.g. web browser**

Inhibit device/user tracking but many organizations want to do the tracking

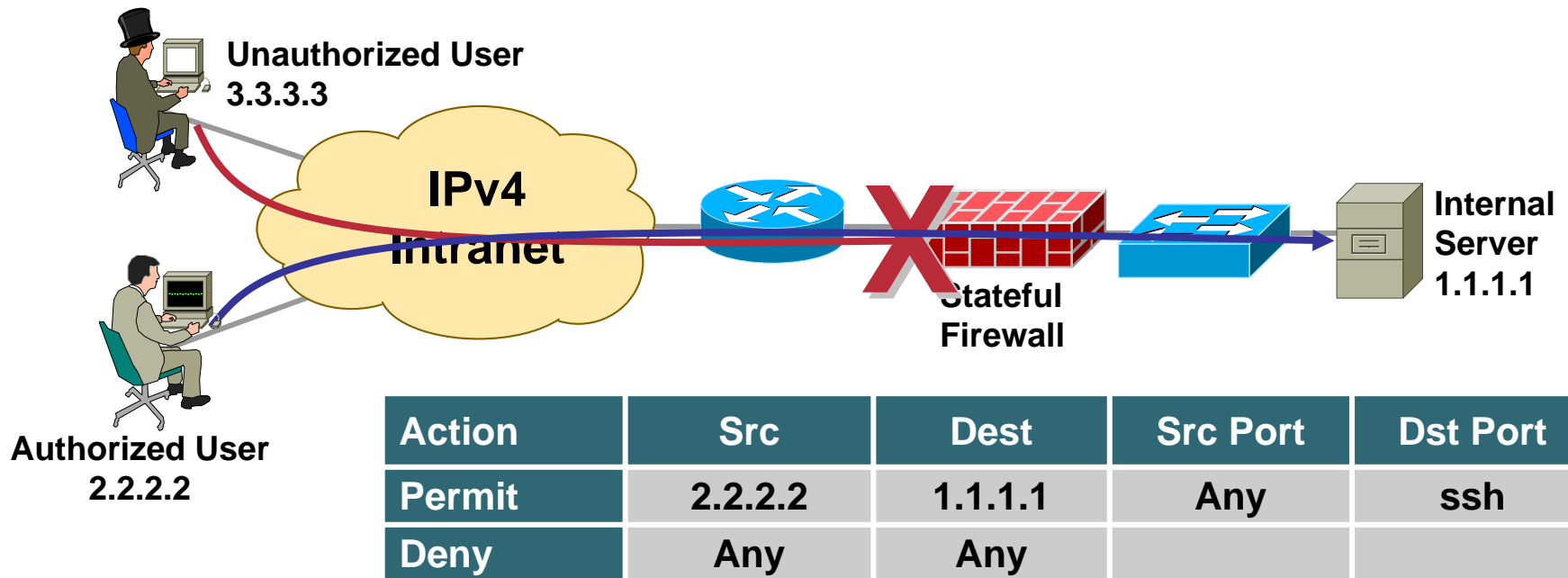
Random 64 bit interface ID, run DAD before using it

Rate of change based on local policy

Recommendation: Use Privacy Extensions for External Communication but Not for Internal Networks (Troubleshooting and Attack Trace Back)

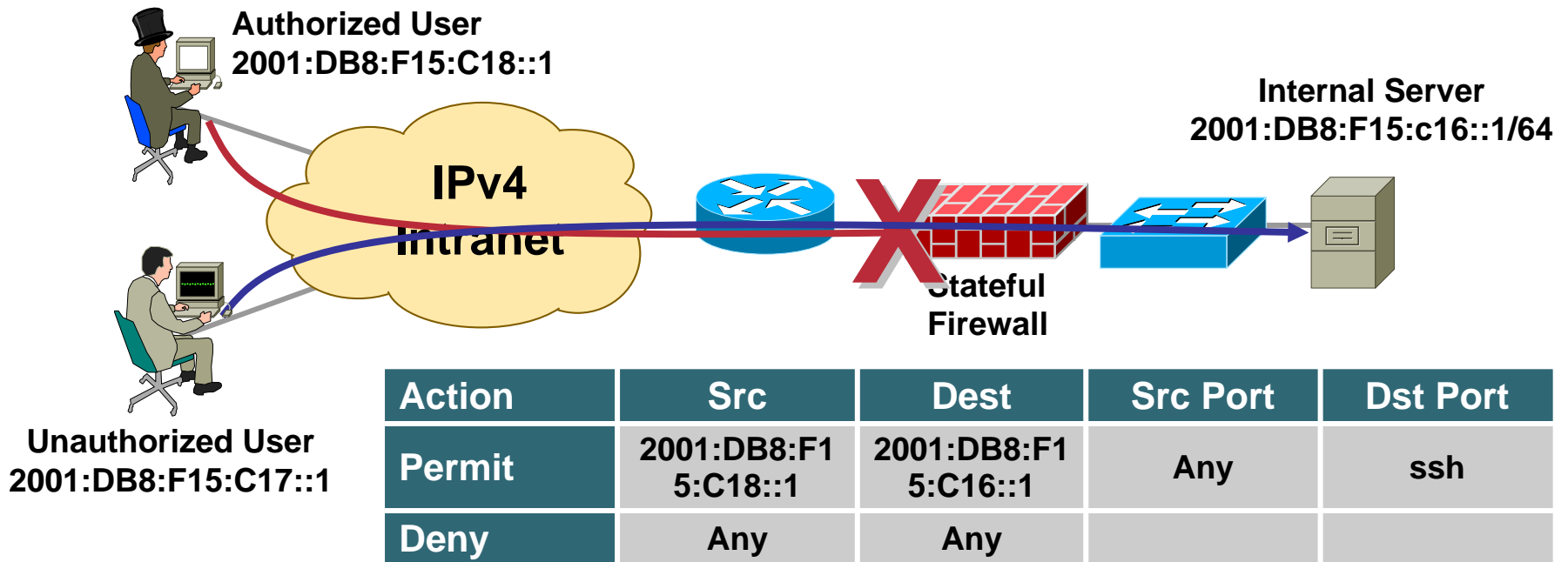
Access Control in IPv4

- Access authorization mainly based on Layer 3 and Layer 4



Access Control in IPv6

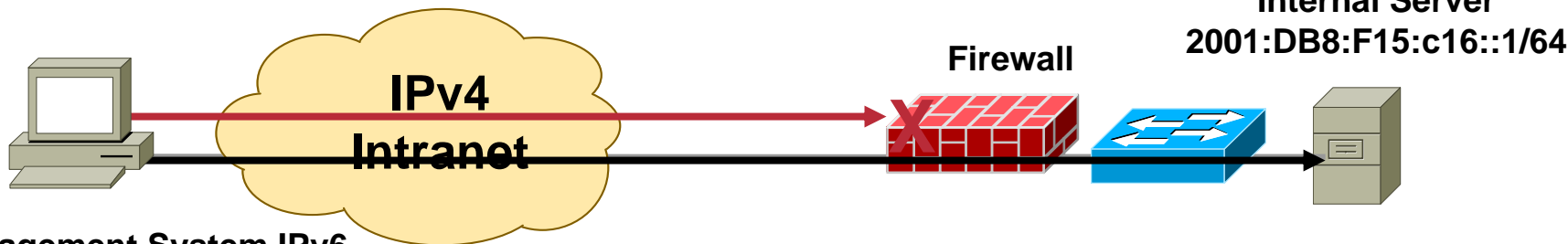
- **Access control in IPv6: Also based on Layer 3 and Layer 4 information**
- **In addition IPv6 has some unique considerations**



Access Control in IPv6 Privacy Extension

- Good to protect the privacy of a host
- But hard to define authorization policy when the Layer 3 information is always changing :-)

Management System New IPv6
Address—2001:DB8:F15:C15::2*



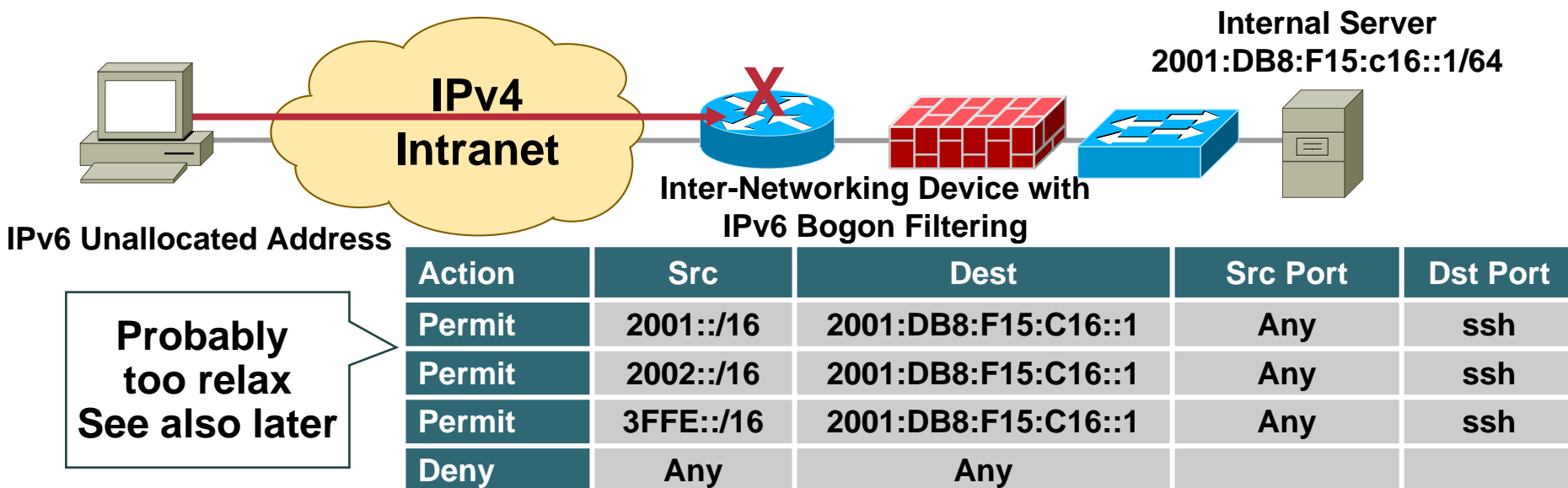
Management System IPv6
Address—2001:DB8:F15:C15::1*

Action	Src	Dest	Src Port	Dst Port
Permit	2001:DB8:F15:C15::1	2001:DB8:F15:c16::1	Any	80
Deny	Any	Any		

*—Not Real
RFC3041 Derived
Addresses

Access Control in IPv6 Bogon Filtering

- In IPv4, it is generally easier to block bogons than it is to permit non-bogons
- In IPv6, a small amount top-level aggregation identifiers (TLAs) have been allocated thus far



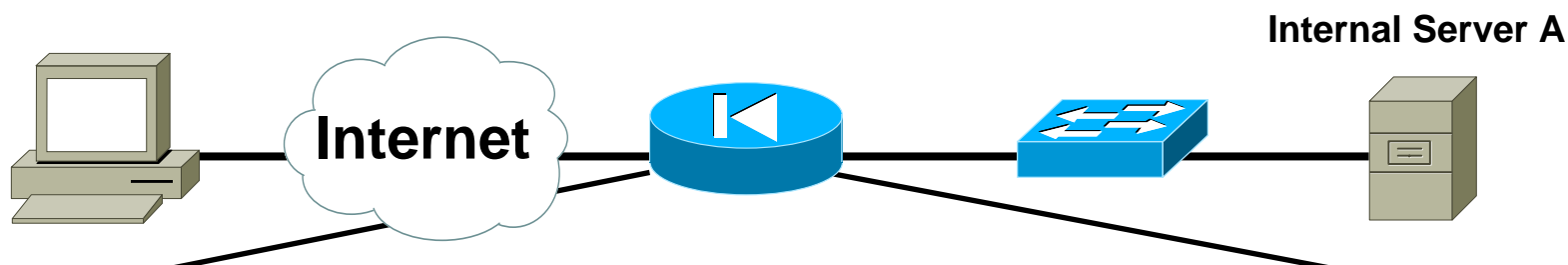
ICMPv4 vs. ICMPv6

- **Significant changes**
- **More relied upon**

ICMP Message Type	ICMPv4	ICMPv6
Connectivity Checks	X	X
Informational/Error Messaging	X	X
Fragmentation Needed Notification	X	X
Address Assignment		X
Address Resolution		X
Multicast Group Management		X
Mobile IPv6 Support		X

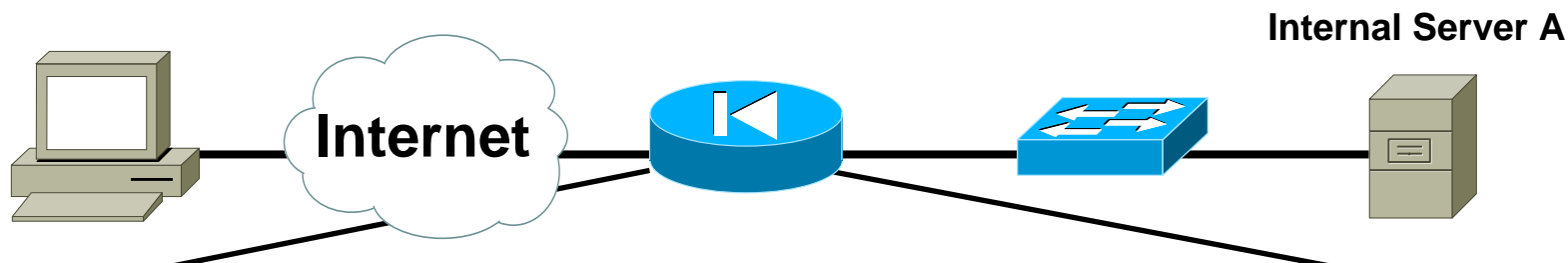
- **=> ICMP policy on firewalls needs to change**

Generic ICMPv4 Border Firewall Policy



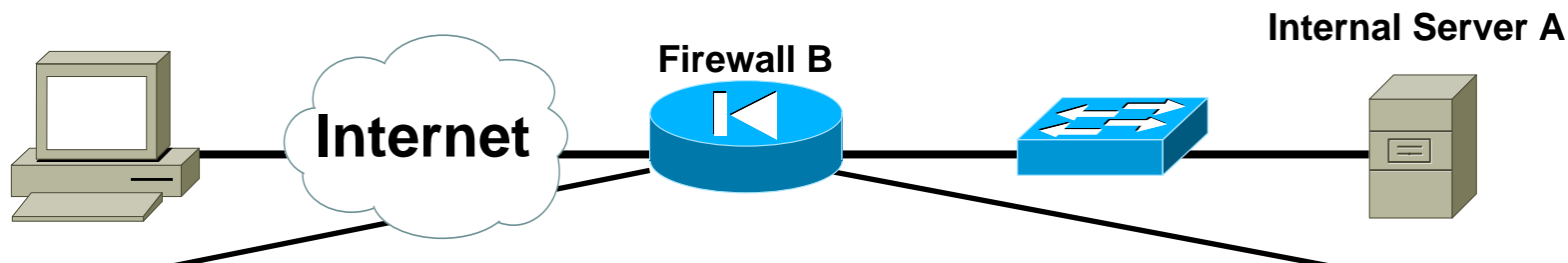
Action	Src	Dst	ICMPv4 Type	ICMPv4 Code	Name
Permit	Any	A	0	0	Echo Reply
Permit	Any	A	8	0	Echo Request
Permit	Any	A	3	0	Dst. Unreachable— Net Unreachable
Permit	Any	A	3	4	Dst. Unreachable— Frag. Needed
Permit	Any	A	11	0	Time Exceeded— TTL Exceeded

Equivalent Comparison ICMPv6 Border Firewall Policy



Action	Src	Dst	ICMPv6 Type	ICMPv6 Code	Name
Permit	Any	A	128	0	Echo Reply
Permit	Any	A	129	0	Echo Request
Permit	Any	A	1	0	No Route to Dst.
Permit	Any	A	2	0	Packet too Big
Permit	Any	A	3	0	Time Exceeded— TTL Exceeded

Potential Additional ICMPv6 Border Firewall Policy



Action	Src	Dst	ICMPv6 Type	ICMPv6 Code	Name
Permit	Any	A	4	0	Parameter Problem
Permit	Any	B	2	0	Packet too Big
Permit	Any	B	130–132	0	Multicast Listener
Permit	Any	B	133/134	0	Neighbor Solicitation and Advertisement
Permit	Any	B	4	0	Parameter Problem

IPv6 Header Manipulation

- Unlimited size of header chain (spec wise) can make filtering difficult
- DoS a possibility with poor IPv6 stack implementations

More boundary conditions to exploit

Can I overrun buffers with a lot of extension headers?

The image shows a Wireshark packet capture of an IPv6 packet. The packet list on the left shows the following structure:

- Frame 1 (423 bytes on wire, 423 bytes captured)
- Raw packet data
- Internet Protocol Version 6
- Hop-by-hop Option Header
- Destination Option Header
- Routing Header, Type 0
- Hop-by-hop Option Header
- Destination Option Header
- Routing Header, Type 0
- Destination Option Header
- Routing Header, Type 0
- Transmission Control Protocol, Src Port: 1024 (1024), Dst Port: bgp (179), Seq: 0, Ack: 0, Len: 51
- Border Gateway Protocol

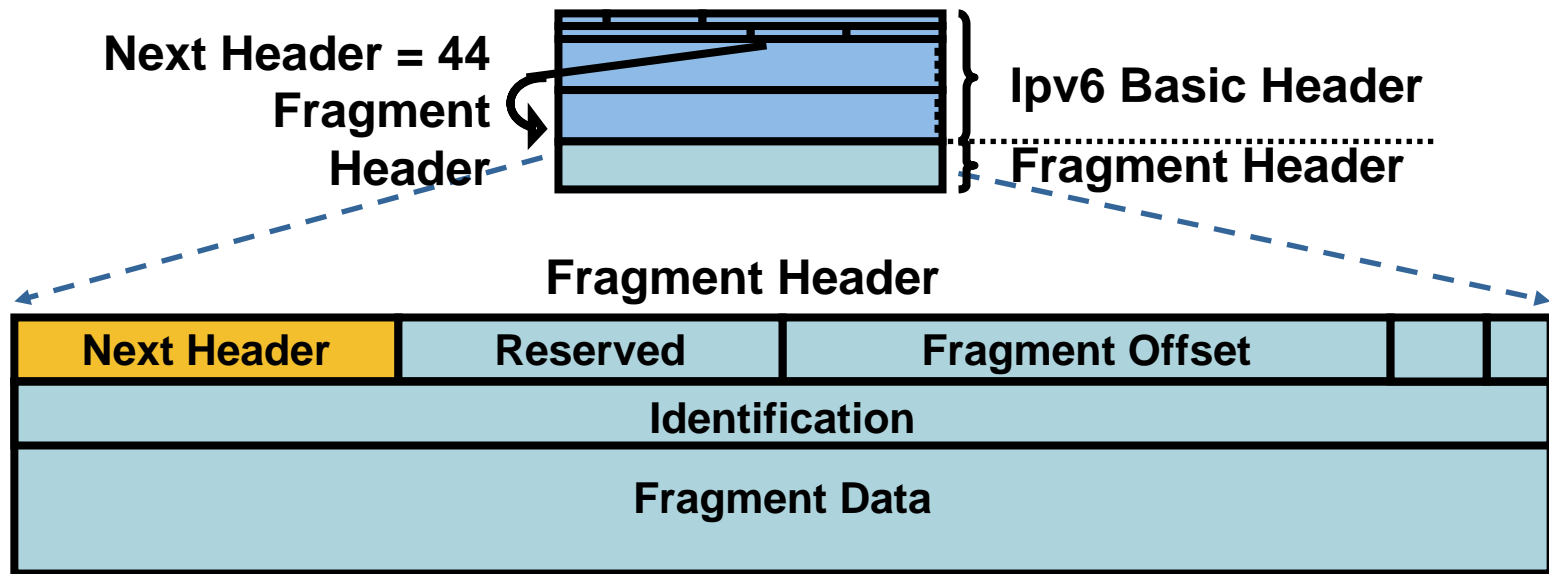
Red circles are drawn around the extension headers in the packet list. Arrows point from these circles to callout boxes on the right:

- A yellow box at the top right states: **Perfectly Valid IPv6 Packet According to the Sniffer**
- A yellow box in the middle right states: **Header Should Only Appear Once** (with an arrow pointing to the first Hop-by-hop Option Header)
- A yellow box below it states: **Destination Header Which Should Occur at Most Twice** (with arrows pointing to the first and second Destination Option Headers)
- A yellow box at the bottom right states: **Destination Options Header Should Be the Last** (with an arrow pointing to the last Destination Option Header)

Fragmentation Used in IPv4 by Attackers

- **Great evasion techniques**
- **Tools like whisker, fragroute, etc.**
- **Makes firewall and network intrusion detection harder**
- **Used mostly in DoSing hosts, but can be used for attacks that compromise the host**

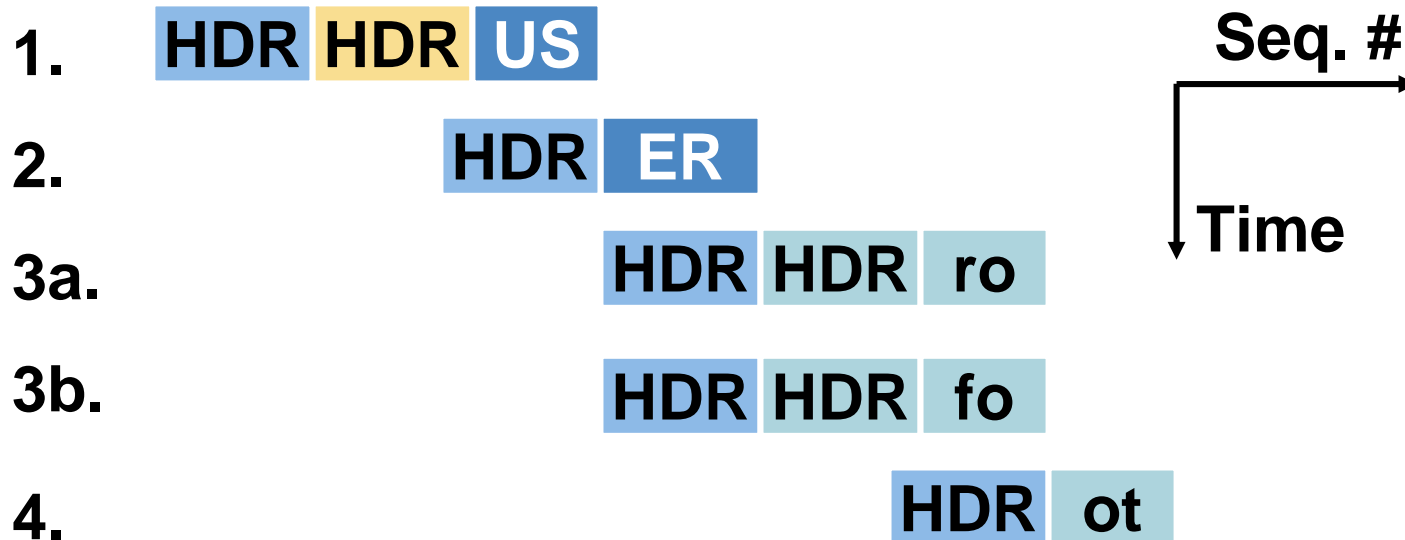
Fragment Header: IPv6



- In IPv6 fragmentation is done **only** by the end system
- Reassembly done by end system like in IPv4
- Attackers can still fragment in intermediate system on purpose
- ==> a great obfuscation tool

IPv6 Fragmentation: Still Need Reassembly in the Firewall and NIDS

Imagine an Attacker Sends:



- Should we consider 3a part of the data stream “USER root”?
- Or is 3b part of the data stream? “USER foot”

If the OS makes a different decision than the monitor: bad

Even worse: different OSs have different protocol interpretations,

If they are overlapping fragments BSD IPv6 drops packet; Linux IPv6 reassembly mimics IPv4 behavior

IPv6 Fragmentation

Issues for Non-Stateful Filtering Devices

- Traverse the next headers before reaching the fragment header to extract the flags and offset
- Then, further NHs before reaching the ULP
- Then check if enough of the upper Layer protocol header is within the first fragment
- This makes matching against the first fragment **non-deterministic**: tcp/udp/icmp might not be there

IPv6 Fragmentation:

Fragment Keyword in IPv6 ACL

- **fragment** keyword matches

Non-initial fragments (same as IPv4)

And the first fragment if the protocol cannot be determined

- **Note: Cisco IOS® also supports a new keyword "undetermined-transport"**

matches any IPv6 packet where the Layer 4 cannot be determined

Header Manipulation and Fragmentation Best Practices

- **Deny IPv6 fragments destined to an internetworking device (DOS vector)**
- **Ensure adequate IPv6 fragmentation filtering capabilities; for example, drop all packets with the routing header if you don't have MIPv6**
- **If really paranoid: drop all fragments with less than 1280 octets (except the last fragment)**

L3-L4 Spoofing in IPv4

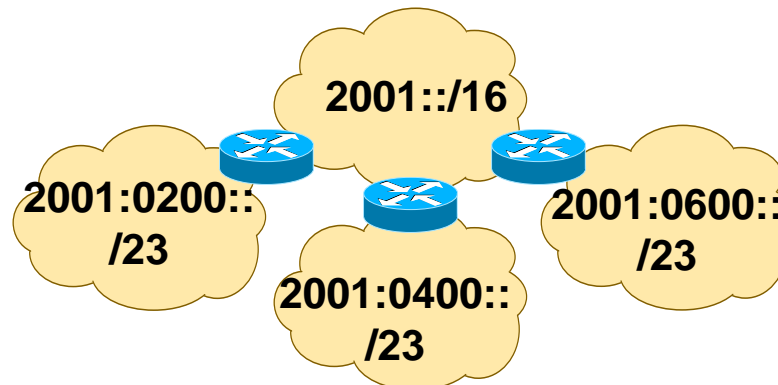
- **L4 spoofing can be done in concert with L3 spoofing to attack systems (most commonly running UDP, i.e. SNMP, Syslog, etc.)**
- **Nearly 50% of the current IPv4 space has not been allocated or is reserved for special use (RFC3330) making it easy to block at network ingress through bogons filtering**

L3-L4 Spoofing in IPv6

IPv6 Address Are Globally Aggregated

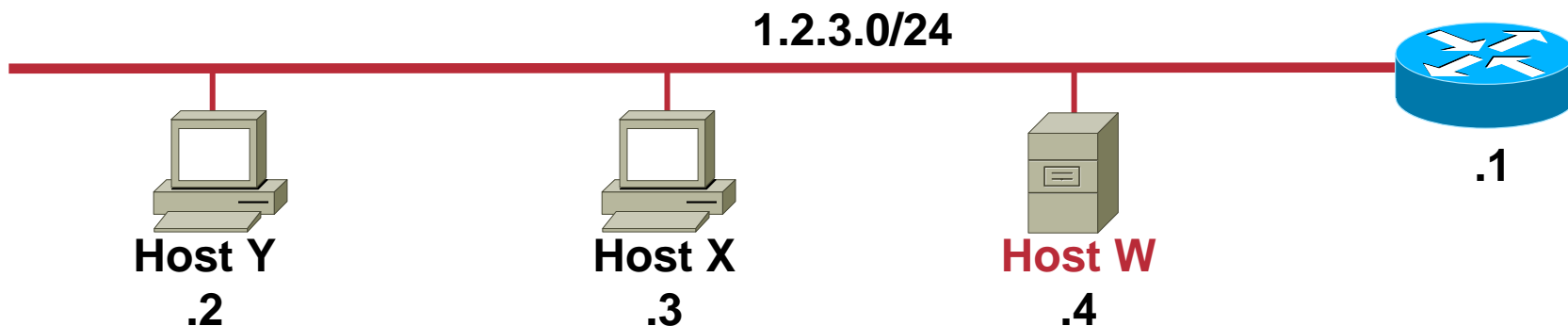
- ==> spoof mitigation at aggregation points easy to deploy
- 2001::/16—IPv6 Production
- 2002::/16—6to4 Tunneling
- 2003::/16—RIPE
- 3FFE::/16—6Bone Testing
- **Unfortunately each subnet (even at the local level) still has a huge range of addresses to spoof**

And now in 2006
2600::/12—ARIN (US DoD)
2A00::/16—RIPE
2400::/16—APNIC



ARP and DHCP Attacks in IPv4

- With ARP misuse host W can claim to be the default gateway and hosts X and Y will route traffic through him; => man in the middle attack



- With DHCP it is similar except the attacker just needs to put a DHCP server on the wire delivering false information (gateways, DNS servers, etc.)

Neighbor Discovery Attacks in IPv6 RFC 3756

- **Redirect attacks**

A malicious node redirects packets away from a legitimate receiver to another node on the link

- **Denial-of-service attacks**

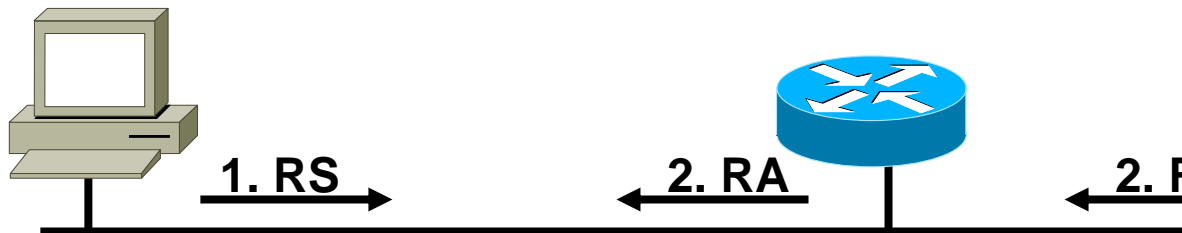
A malicious node prevents communication between the node under attack and other nodes

- **Flooding denial-of-service attacks**

A malicious node redirects other hosts' traffic to a victim node creating a flood of bogus traffic at the victim host

Stateless Autoconfiguration

Router Solicitation Are Sent By Booting Nodes to Request Router Advertisements for Configuring the Interfaces



ICMP w/o any authentication
Gives Exactly Same Level of Security as ARP For IPv4 (None)
Bootstrap Security Problem Just Like IPv4

1. RS:

ICMP Type = 133

Src = ::

Dst = All-Routers multicast Address

query= please send RA

2. RA:

ICMP Type = 134

Src = Router Link-local Address

Dst = All-nodes multicast address

Data= options, prefix, lifetime, **autoconfig** flag

Neighbor Discovery: Neighbor Solicitation



ICMP type = 135

Src = A

Dst = Solicited-node multicast of B

Data = link-layer address of A

Query = what is your link address?

**Security Mechanisms
Built into Discovery
Protocol = None
Another Bootstrap
Security Problem**



ICMP type = 136

Src = B

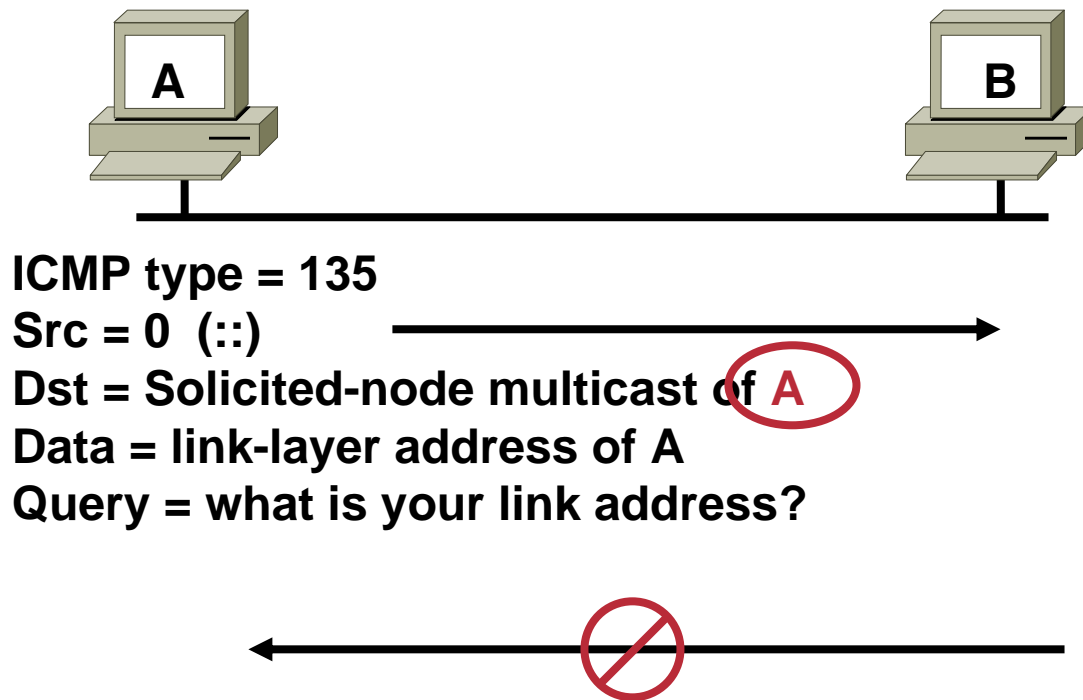
Dst = A

Data = link-layer address of B

**A and B Can Now Exchange
Packets on This Link**

DAD (Duplicate Address Detection)

Duplicate Address Detection (DAD) Uses Neighbor Solicitation to Verify the Existence of an Address to be Configured

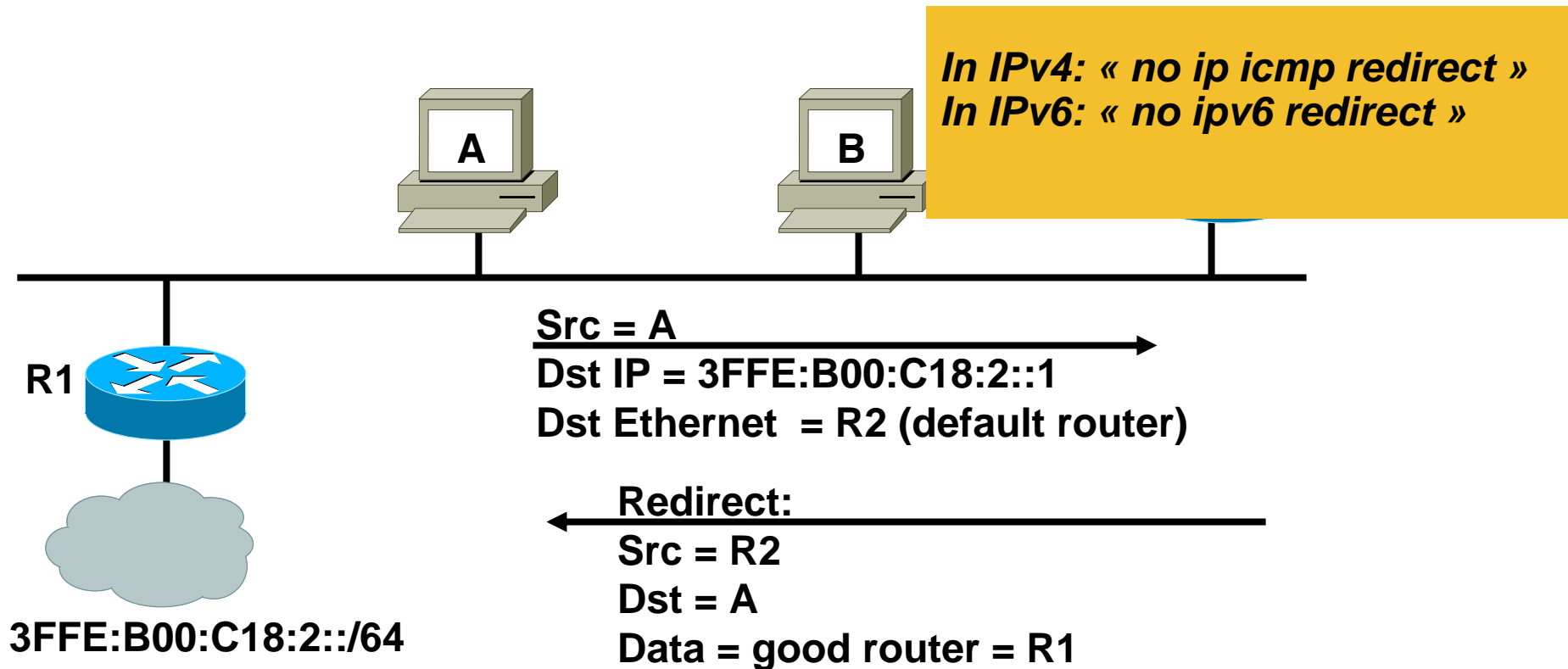


From RFC 2462:
« If a Duplicate @ Is Discovered... the Address Cannot Be Assigned to the Interface »

⇔ **What If: Use MAC@ of the Node You Want to Dos and Fabricate Its IPv6 @**

Neighbor Discovery: Spoofed Redirect

Redirect is Used by a Router to Signal the Re-Route of a Packet to a Better Router



Secure Neighbor Discovery (SEND)

RFC 3971

- **Certification paths**

Anchored on trusted parties, expected to certify the authority of the routers

- **Cryptographically Generated Addresses (CGA)**

IPv6 addresses whose the interface identifier is cryptographically generated

- **RSA signature option**

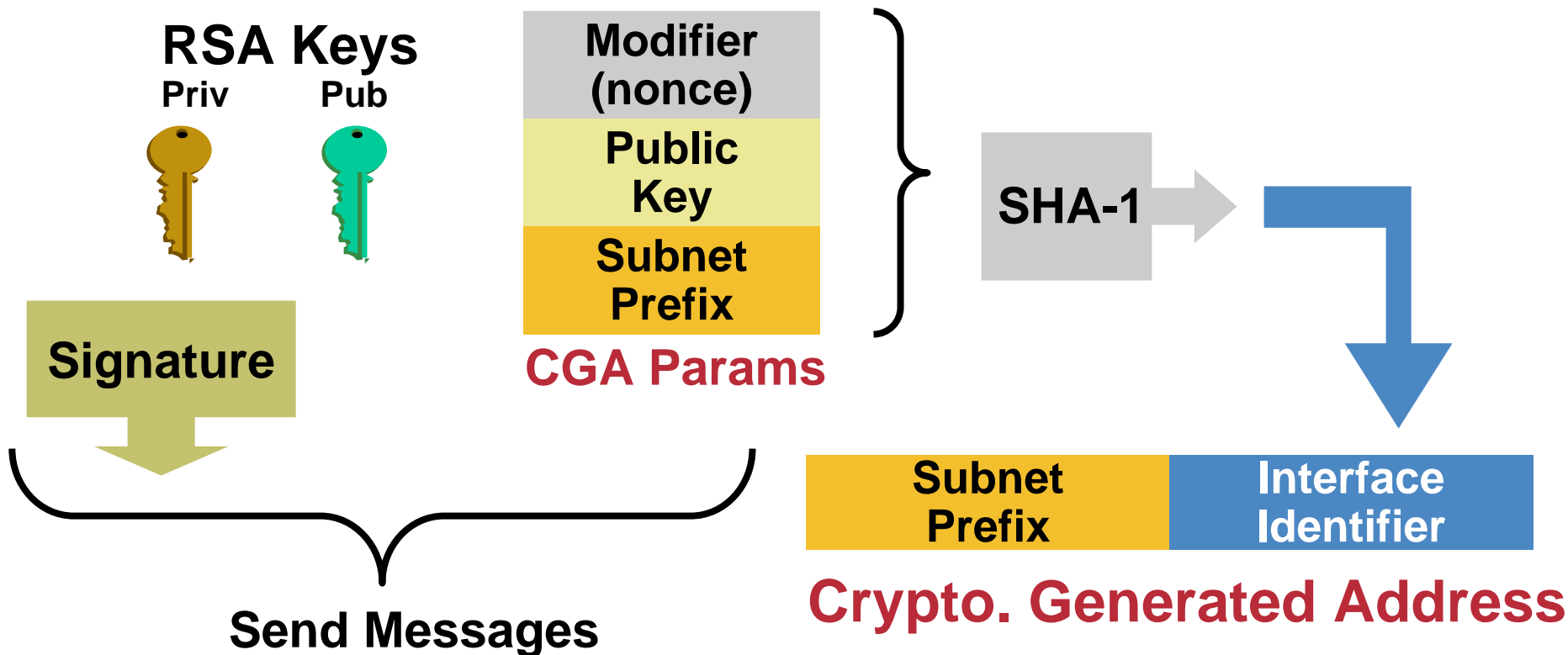
Protect all all messages relating to neighbor and router discovery

- **Timestamp and nonce ND options**

Prevent replay attacks

Cryptographically Generated Addresses CGA RFC 3972 (Simplified)

- Each device has a RSA key pair (no need for cert)
- Ultra light check for validity
- Prevent spoofing a valid CGA address



Secure Neighbor Discovery and CGA Caveats

- **Private/public key pair on all devices for CGA**
- **Overhead introduced**

**Routers have to do many public/private key calculation
(some may be done in advance of use)**

- **Available: Linux**
- **Coming in Microsoft Vista SP1**
- **Future implementation: Cisco IOS**

DHCPv6 Threats

- **Note: use of DHCP is announced in Router Advertisements**
- **Rogue devices on the network giving misleading information or consuming resources (DoS)**
 - Rogue DHCPv6 client and servers on the network (same threat as IPv4)**
 - Rogue DHCPv6 servers on the site local multicast address (FF05::1:3) (new threat in IPv6)**
- **Tampering of communication between DHCPv6 relays and servers**

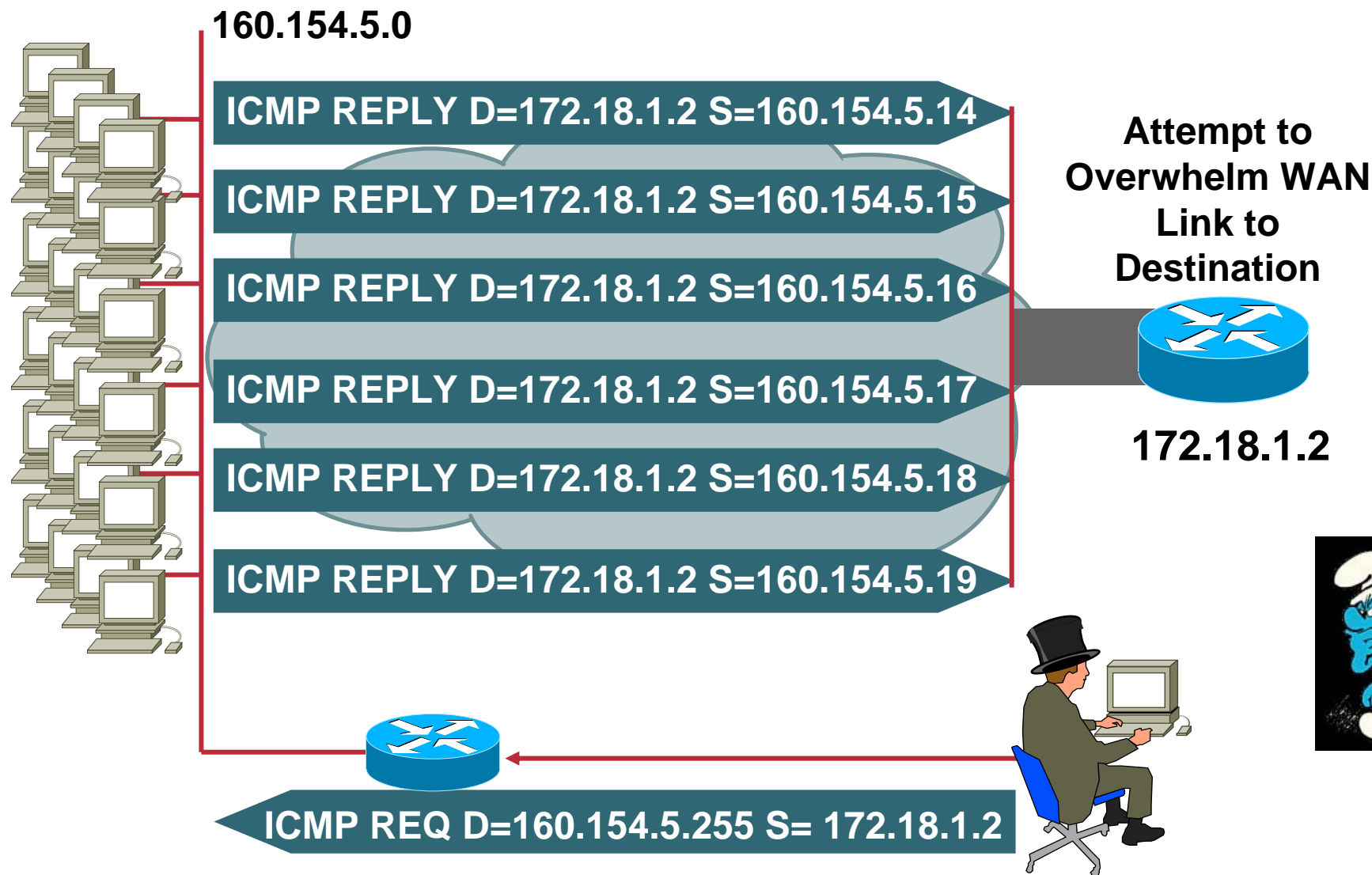
DHCPv6 Threat Mitigation

- **Rogue clients and servers can be mitigated by using the authentication option in DHCPv6**

There are not many DHCPv6 client or server implementations using this today.

- **For paranoid: protect the relay to server communications with IPsec (similar to IPv4)**

IPv4 Broadcast Amplification: Smurf



IPv6 and Broadcasts

- **There are no broadcast addresses in IPv6**
- **Broadcast address functionality is replaced with the appropriate link local multicast address**

Link Local All Nodes Multicast—FF02::1

Link Local All Routers Multicast—FF02::2

IPv6 and Other Amplification Vectors

- **Specific mention is made in ICMPv6 RFC that no ICMP error message should be generated in response to a packet with a multicast destination address**
- **The exceptions are the packet too big message and the parameter problem ICMP messages**
- **RFC 2463 Section 2.4 (e.2)**

**Implement Ingress Filtering of Packets With
IPv6 Multicast Source Addresses
IPv6 Mcast Dest Address and Above ICMP Packet Types**

Preventing IPv6 Routing Attacks

Protocol Authentication

- **BGP, ISIS, EIGRP no change:**

An MD5 authentication of the routing update

- **OSPFv3 has changed and pulled MD5 authentication from the protocol and instead is supposed to rely on transport mode IPsec**

- **RIPng also relies on IPsec**

- **IPv6 routing attack best practices**

Use traditional authentication mechanisms on BGP and IS-IS

Use IPsec to secure protocols such as OSPFv3 and RIPng

Viruses and Worms in IPv6

- **Pure viruses don't change in IPv6 but hybrid and pure worms do**

Hybrids and pure worms today rely in Internet scanning to infect other hosts, this isn't feasible as shown earlier in this presentation

At one million packets per second on a IPv6 subnet with 10,000 hosts it would take over 28 years to find the first host to infect

- **Worm developers will adapt to IPv6 but pure random scanning worms will be much more problematic for the attacker; best practices around worm detection and mitigation from IPv4 remain**

IPv6 Attacks with Strong IPv4 Similarities

- **Sniffing**

Without IPsec, IPv6 is no more or less likely to fall victim to a sniffing attack than IPv4

- **Application layer attacks**

Even with IPsec, the majority of vulnerabilities on the Internet today are at the application layer, something that IPsec will do nothing to prevent

- **Rogue devices**

Rogue devices will be as easy to insert into an IPv6 network as in IPv4

- **Man-in-the-Middle Attacks (MITM)**

Without IPsec, any attacks utilizing MITM will have the same likelihood in IPv6 as in IPv4

- **Flooding**

Flooding attacks are identical between IPv4 and IPv6

By the Way: It Is Real ☹️

IPv6 Hacking Tools

Let the Games Begin

- **Sniffers/packet capture**

- Snort
- TCPdump
- Sun Solaris snoop
- COLD
- Ethereal
- Analyzer
- Windump
- WinPcap
- NetPeek
- Sniffer Pro

- **Worms**

- Slapper

- **Advisories/field notices**

- <http://www.cisco.com/warp/public/707/cisco-sa-20050126-ipv6.shtml>

- <http://www.kb.cert.org/vuls/id/658859>

- **Scanners**

- IPv6 security scanner
- Halfscan6
- Nmap
- Strobe
- Netcat

- **DoS Tools**

- 6tunneldos
- 4to6ddos
- Imps6-tools

- **Packet forgers**

- SendIP
- Packit
- Spak6

- **Complete tool**

- <http://www.thc.org/thc-ipv6/>

IPv6 Security Best Practice



Wrap Up: Candidate Best Practices

- **Implement privacy extensions carefully**
- **Filter internal-use IPv6 addresses at the enterprise border routers**
- **Filter unneeded services at the firewall**
- **Selectively filter ICMP**
- **Maintain host and application security**
- **Determine what extension headers will be allowed through the access control device**
- **Determine which ICMPv6 messages are required**
- **Deny IPv6 fragments destined to an internetworking device when possible**
- **Ensure adequate IPv6 fragmentation filtering capabilities**
- **Drop all fragments with less than 1280 octets (except the last one)**

Wrap Up: Candidate Best Practices (Cont.)

- **Implement RFC 2827-like filtering and encourage your ISP to do the same**
- **Document procedures for last-hop traceback**
- **Use cryptographic protections where critical**
- **Use static neighbor entries for critical systems**
- **Implement ingress filtering of packets with IPv6 multicast source addresses**
- **Use traditional authentication mechanisms on BGP and IS-IS**
- **Use IPsec to secure protocols such as OSPFv3 and RIPng**
- **Use IPv6 hop limits to protect network devices**
- **Use static tunneling rather than dynamic tunneling**
- **Implement outbound filtering on firewall devices to allow only authorized tunneling endpoints**

Specific IPv6 Issues



IPv6 Transition Mechanism Challenges

- **16+ methods, possibly in combination**

IP spoofing

- **Dual stack**

Consider security for both protocols

Cross v4/v6 abuse

Resiliency (shared resources)

- **Tunnels**

Bypass firewalls (protocol 41)

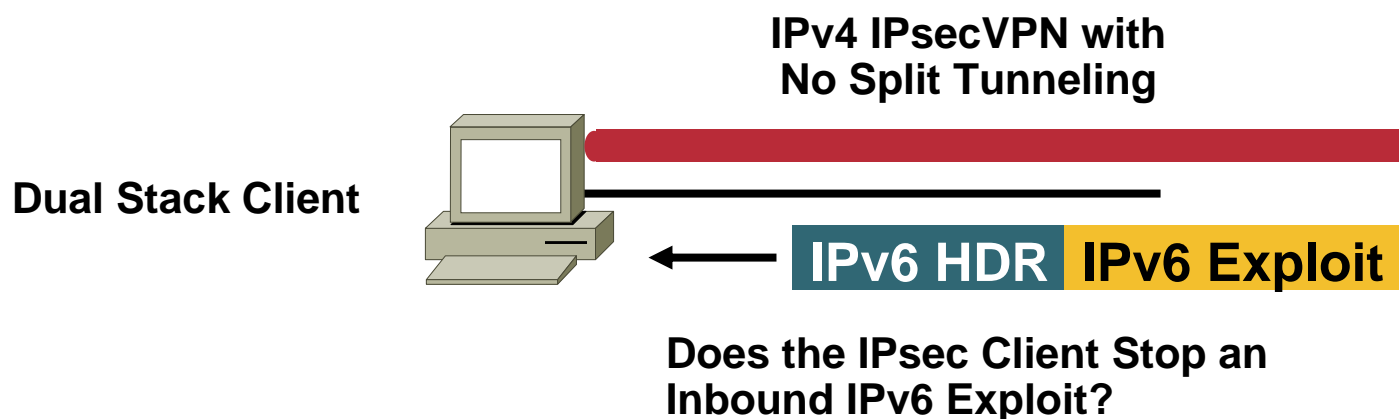
IPv6 Dual Stack Host Considerations

- **Host security on a dual-stack device**

Applications can be subject to attack on both IPv6 and IPv4

- **Host security controls should block and inspect traffic from both IP versions**

Host intrusion prevention, personal firewalls, VPN clients, etc.



IPv6 Tunneling Summary

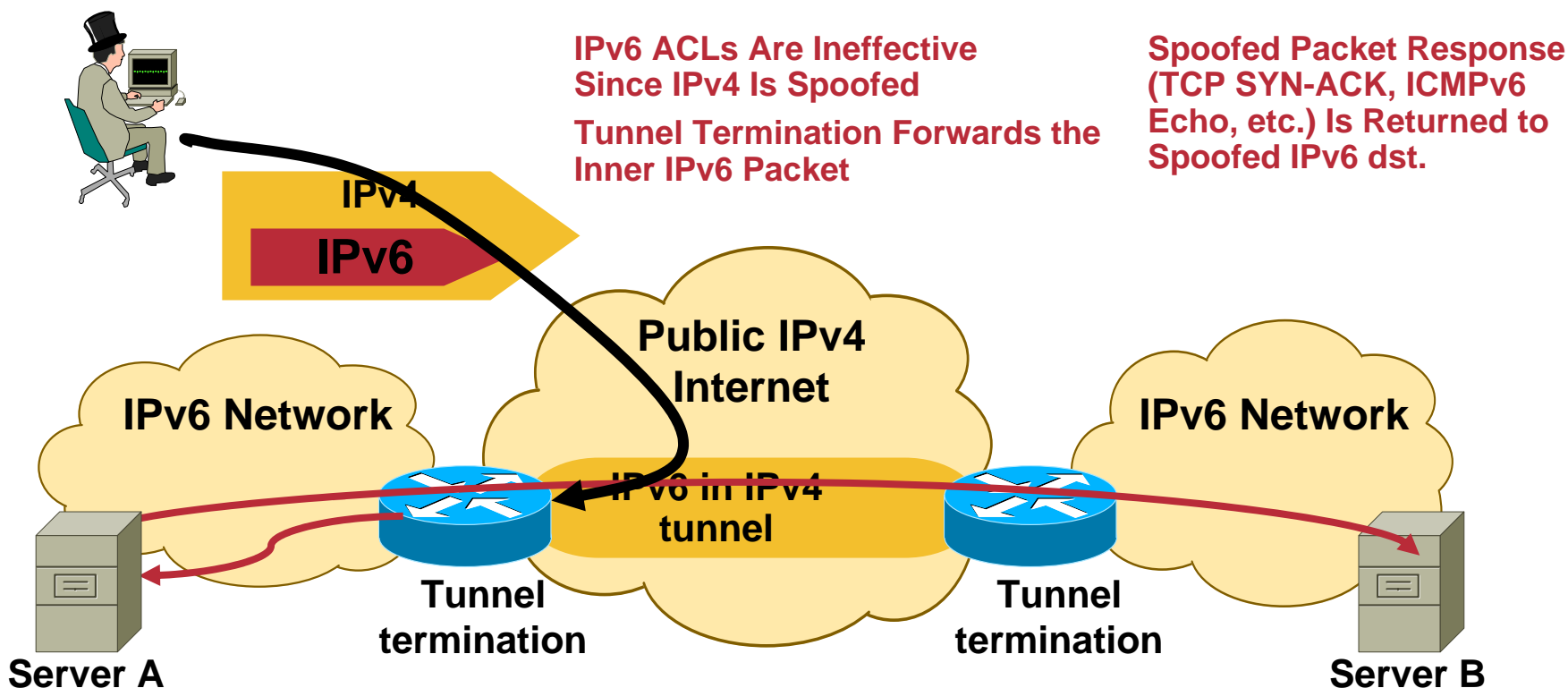
- RFC 1933/2893 configured and automatic tunnels
- RFC 2401 IPsec tunnel
- RFC 2473 IPv6 generic packet tunnel
- RFC 2529 6over4 tunnel
- RFC 3056 6to4 tunnel
- ISATAP tunnel
- MobileIPv6 (uses RFC2473)
- Teredo tunnels

- Only allow authorized endpoints to establish tunnels
- Static tunnels are deemed as “more secure,” but less scalable
- Automatic tunneling mechanisms are susceptible to packet forgery and DoS attacks
- Automatic tunneling mechanisms are susceptible to packet forgery and DoS attacks
- These tools have the **same risk** as IPv4, just new avenues of exploitation
- Automatic IPv6 over IPv4 tunnels could be secured by IPv4 IPsec

L3-L4 Spoofing in IPv6

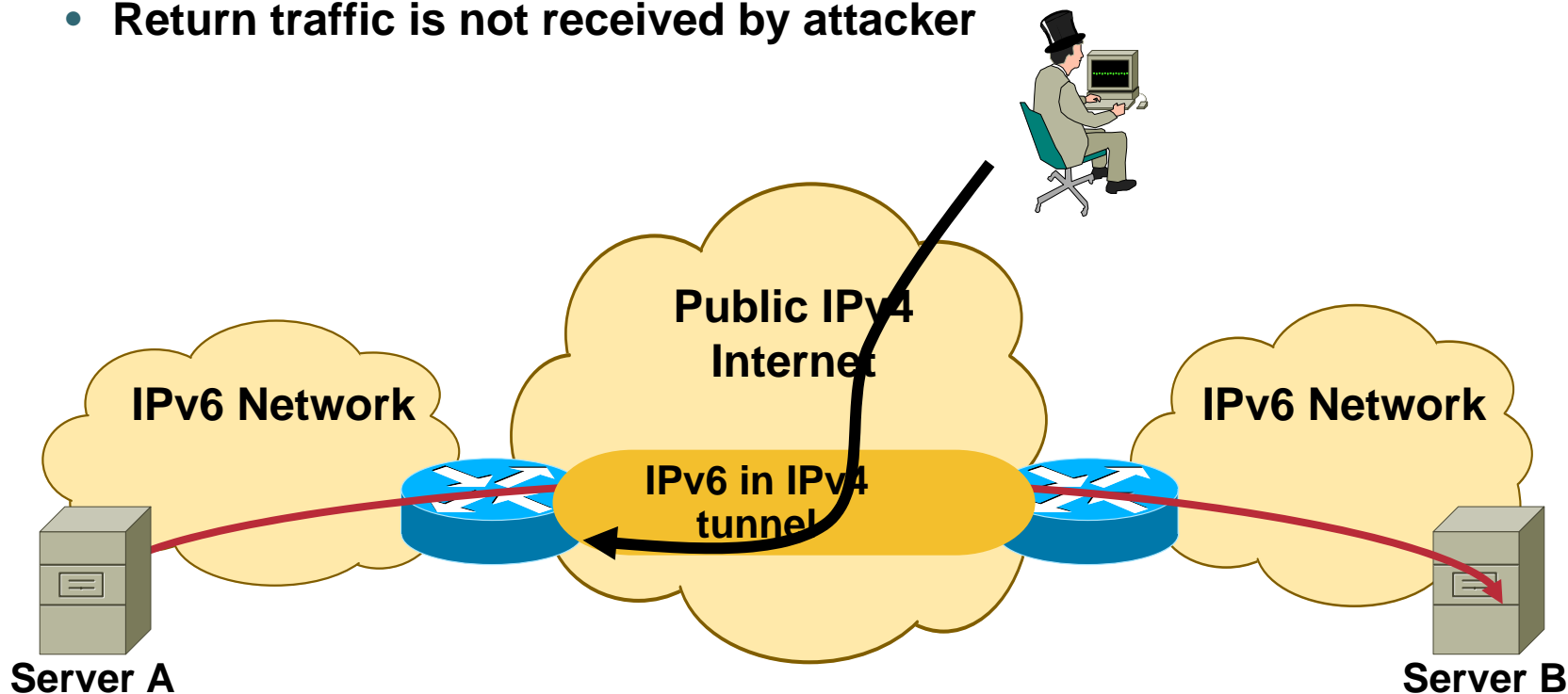
When Using IPv6 over IPv4 Tunnels

- Most IPv4/IPv6 transition mechanisms have no authentication built in
- => an IPv4 attacker can inject traffic if spoofing on IPv4 and IPv6 addresses



L3-L4 Spoofing in IPv6 Via Tunnels

- Harm is limited
- 1:1 ratio of packets—no amplification attack
- There is a chokepoint against DoS
- Return traffic is not received by attacker



Transition Threats

- **ISATAP threats**

Unauthorized tunnels—firewall bypass (protocol 41)

ISATAP looks like a Layer 2 network to ALL ISATAP hosts in the enterprise

This has implications on network segmentation and network discovery

No authentication in ISATAP—rogue routers are possible

Host security needs IPv6 support

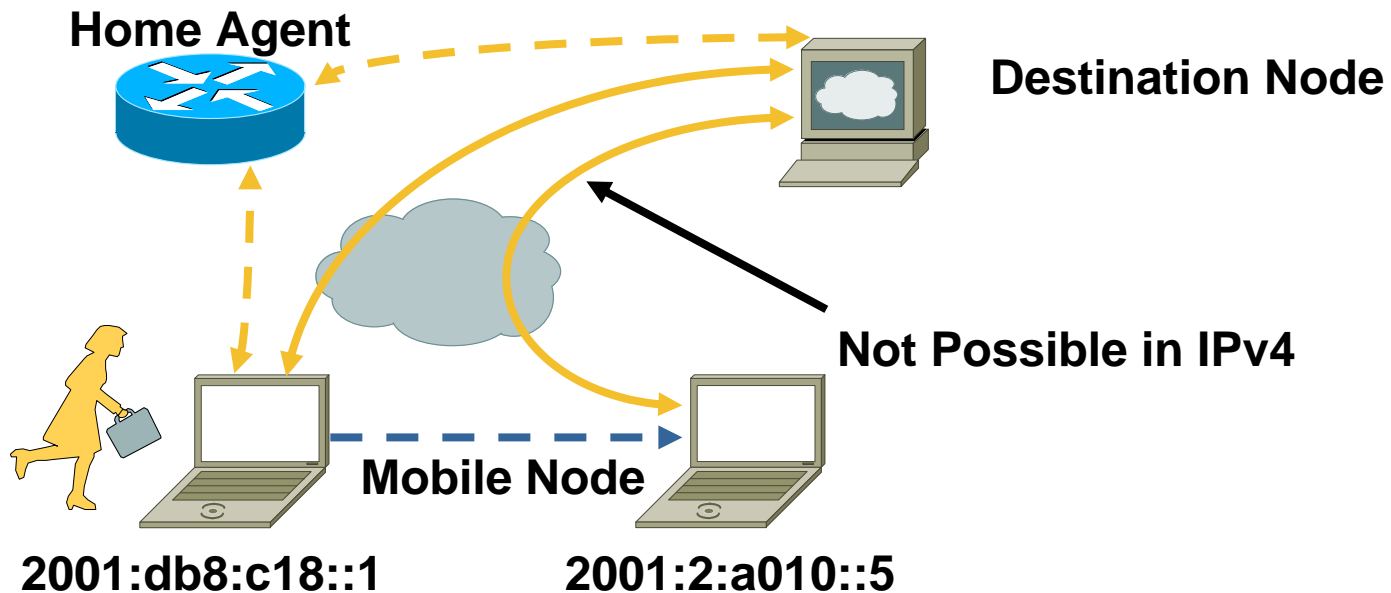
- **Teredo threats—IPv6 over UDP (port 3544)**

Unauthorized tunnels—firewall bypass

Rogue relays/servers can be used for DoS; a possible for client to server communications

Host security needs IPv6 support

IP Mobility



Mobility Means:

- Mobile devices are fully supported while moving
- Built-in on IPv6
 - Any node can use it
- Efficient routing means performance for end-users
- **Filtering challenges**

Mobile IPv6 Security Features Overview

- **Protection of binding updates both to home agents and correspondent nodes**

IPsec,

Or binding authorization data option through the **return routability** procedure

- **Protection of mobile prefix discovery**

Through the use of IPsec extension headers

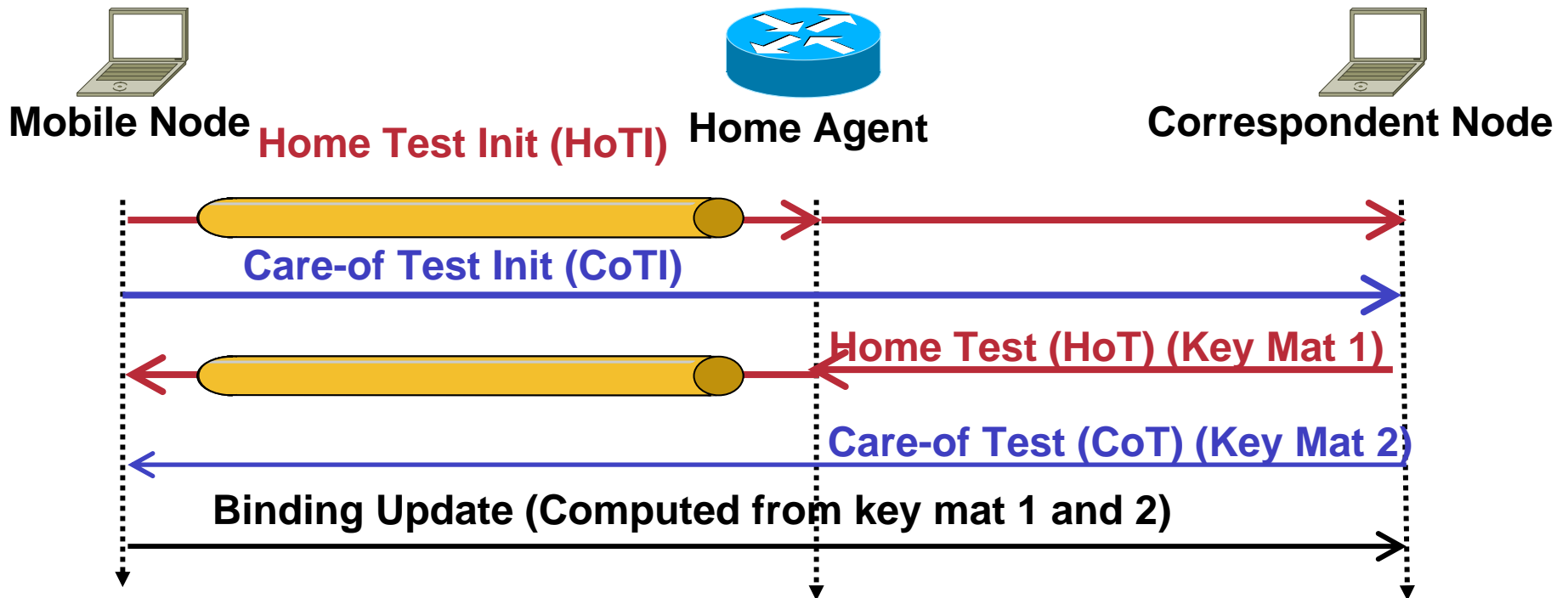
- **Protection of data packets transport**

Home address destination option and **type two routing** header specified in a manner which restricts their use in attacks

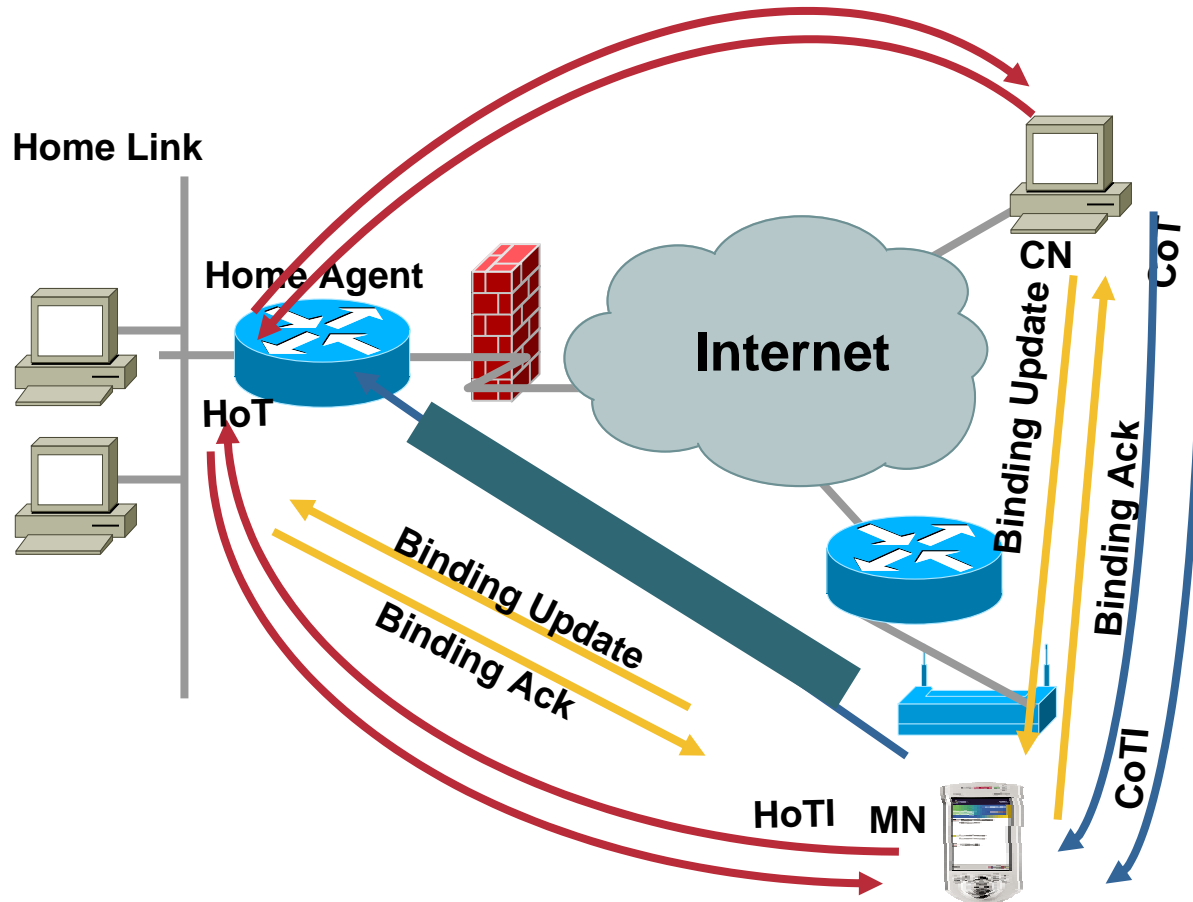
Mobile IPv6 Security

Return Routability Test

- Provides reasonable assurance that the MN is addressable at its claimed CoA and at its HoA
- Test whether packets addressed to the two claimed addresses are routed back to the MN



Mobile IPv6 Global Picture



- **Correspondent Node**

Arbitrary: No
Preexisting Security
Association

- **Return Routability Test**

Verifies the
collocation of the
CoA and the home
address

Assumes better
security association
between HA and MN

Scalable and
stateless

- **Reverse Tunnel**

Secured by IPsec

Requires a
preexisting Security
Association

MIPv6 Security Protections

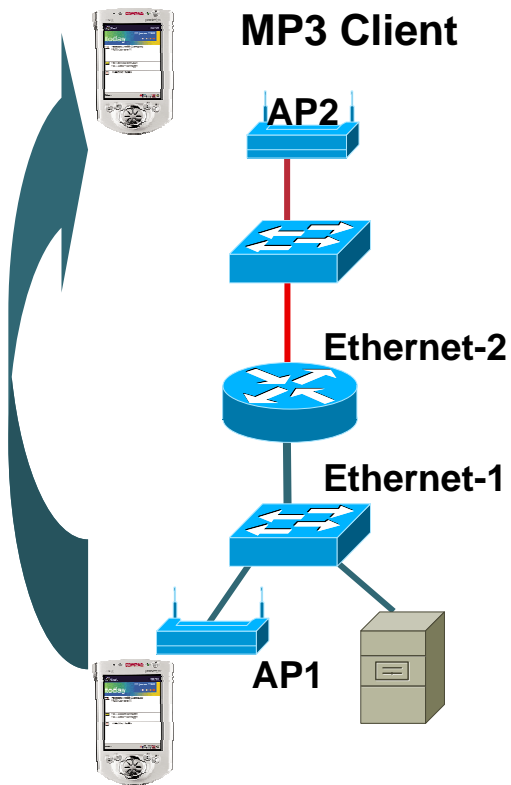
- BU/BA to HA **must** be secured through IPsec
- MN and HA **should** use an IPsec SA to protect the integrity and authenticity of the mobile prefix solicitations and advertisements
- Payload packets exchanged with MN can be follow the same protection policy as other IPv6 hosts
- Specific security measures are defined to protect the specificity of MIPv6

Home address destination option

Type 2 Routing header

Tunnelling headers

Mobile IPv6 ACL



Router# (config-if) ipv6 mobile home-agent access <acl>

- **Binding update filter: all received binding updates are filtered**
- **This feature may be used to deny home agent services to mobile nodes that have roamed to particular sub-networks**

When the filter blocks a binding update, a binding acknowledgement is returned with error status “administratively prohibited”

Enforcing a Security Policy




Cisco IOS IPv6 Access Control Lists

- **Can filter traffic based on source and destination address**
- **Can filter traffic inbound or outbound to a specific interface**
- **Implicit "deny all" at the end of access list**
- **Very much like in IPv4**

Cisco IOS IPv6 Access Control Lists

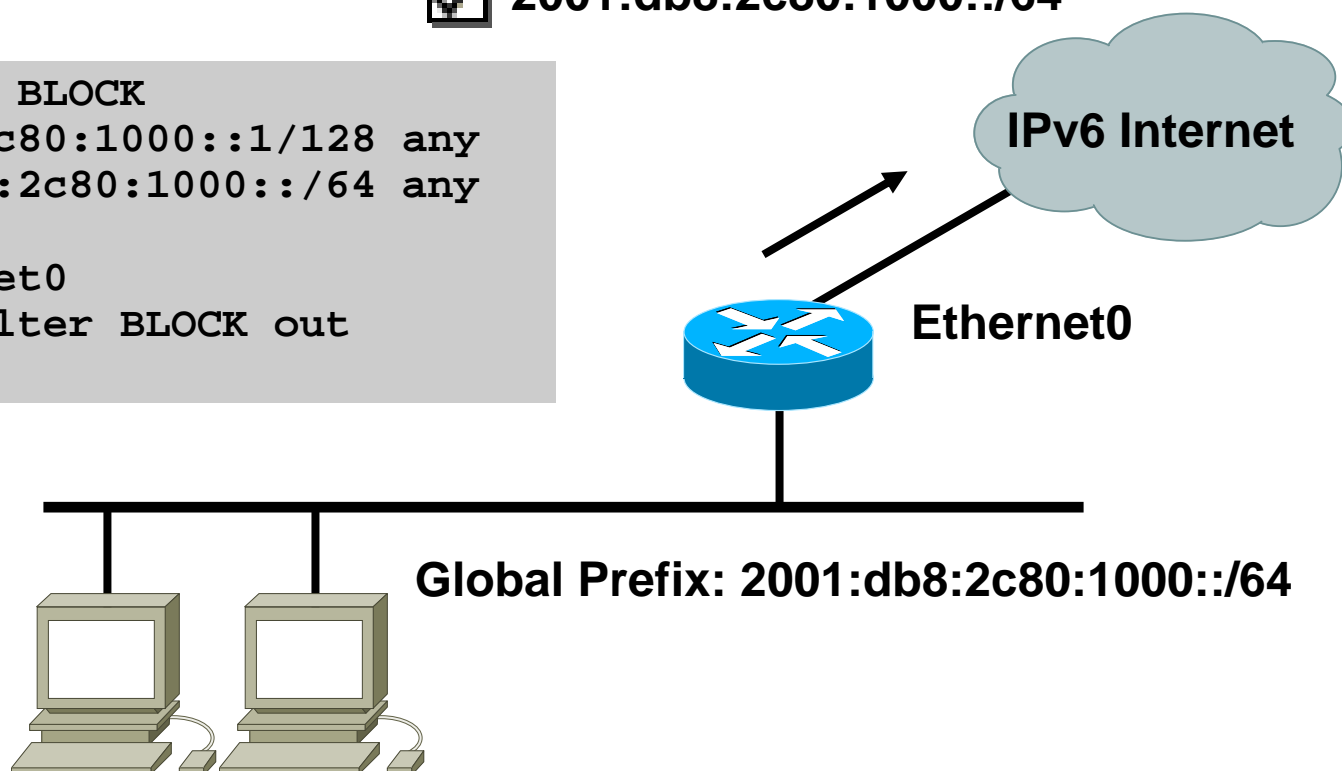
Filtering Outgoing Traffic from One Specific Source Address

 2001:db8:2c80:1000::1

 2001:db8:2c80:1000::/64

```
ipv6 access-list BLOCK
deny 2001:db8:2c80:1000::1/128 any
permit 2001:db8:2c80:1000::/64 any
```

```
interface Ethernet0
ipv6 traffic-filter BLOCK out
```



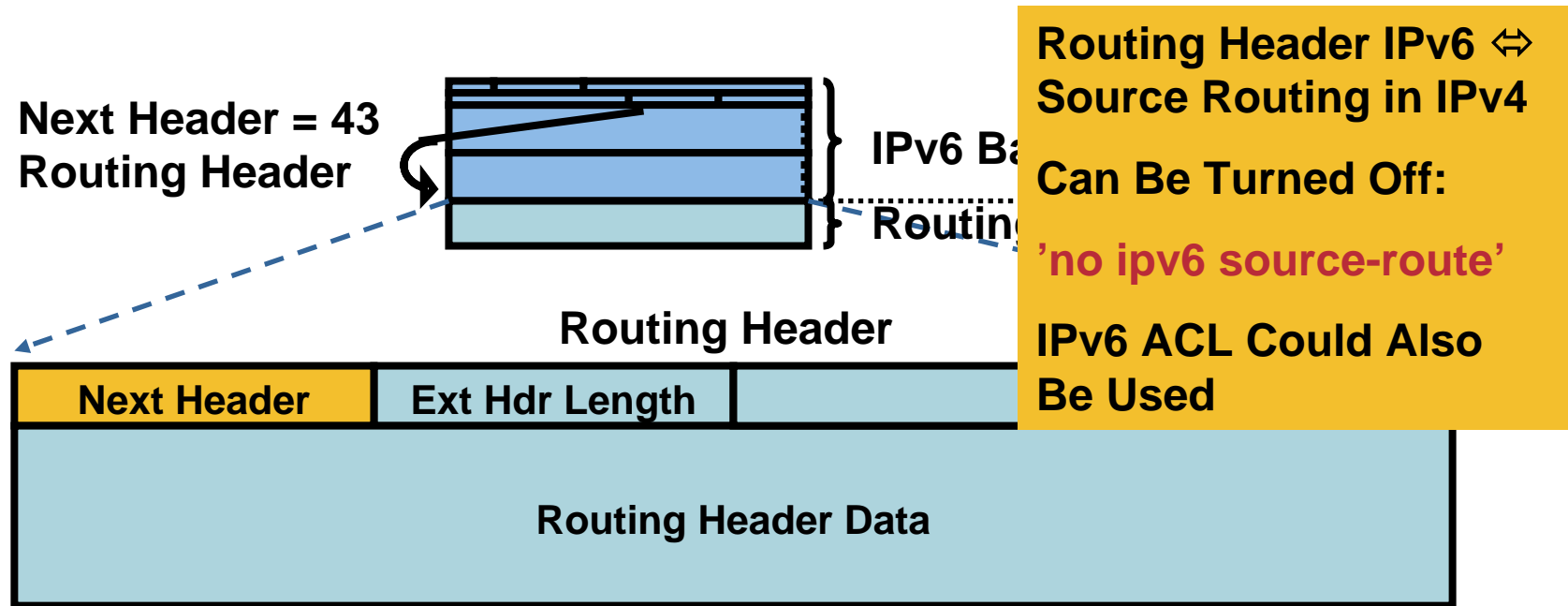
Filtering Extension Headers

- IPv6 headers and optional extensions need to be scanned to access the upper layer protocols (UPL)
- May require searching through several extensions headers
- **Important:** a router must be able to filter both option header and L4 at the same time

IPv6 Routing Header

Routing Header Is:

- An extension header
- Processed by the listed intermediate routers



Issues with Routing Header

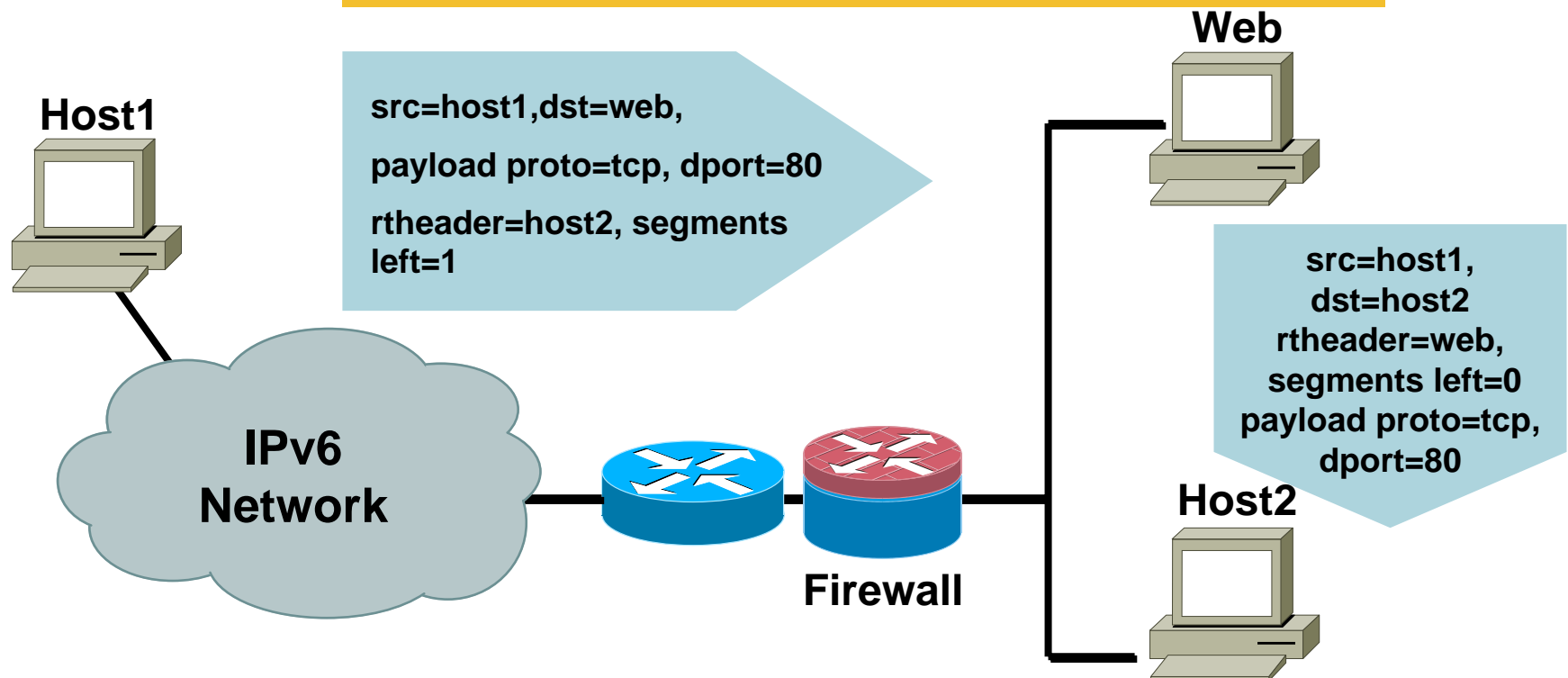
- **Could be used as a rebound/relay to the victim**
- **Because destination address is replaced at every routing header processing point, it's difficult to perform traffic filtering based on destination addresses**
- **<http://www.ietf.org/internet-drafts/draft-savola-ipv6-rh-ha-security-03.txt>**

Routing Header: Traffic Reflector

Rule on the Firewall

Allow proto tcp from any to webserver port 80

Deny proto tcp from any to any



IPv6 Extended Access Control Lists

- **Upper layers : ICMP, TCP, UDP, SCTP, any value**
- **ICMPv6 code and type**
- **TCP SYN, ACK, FIN, PUSH, URG, RST**
- **L4 port numbers**
- **Traffic class (only six bits/8) = DSCP**
- **Flow label (0-0xFFFFF)**
- **IPv6 header options**
 - Fragments**
 - Routing header type**
 - Destination header type**

IPv6 ACL Implicit Rules

Implicit Permit for Enable Neighbor Discovery

- The following implicit rules exist at the end of each IPv6 ACL to allow ICMPv6 neighbor discovery:

```
permit icmp any any nd-na  
permit icmp any any nd-ns  
deny ipv6 any any
```

- Be careful when adding « `deny ipv6 any any log` » at the end

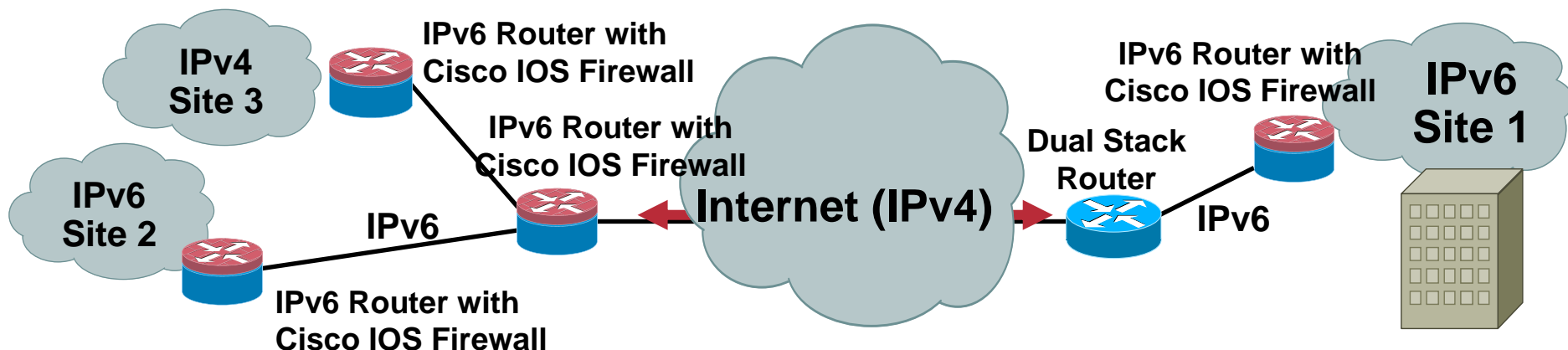
IPv6 ACL to Protect VTY

```
ipv6 access-list VTY
  permit ipv6 2001:db8:0:1::/64 any

line vty 0 4
  ipv6 access-class VTY in
```

Cisco IOS Firewall IPv6 Support

- Stateful protocol inspection (anomaly detection) of IPv6 fragmented packets, TCP, UDP, ICMP and FTP traffic
- Stateful inspection and translation services of IPv4/IPv6 packets
- IPv6 DoS attack mitigation
- IPv4/v6 coexistence, no need for new hardware, just software
- Recognizes IPv6 extension header information such as routing header, hop-by-hop options header, fragment header, etc



Control Plane Policing for IPv6

Protecting the Router CPU

- Against DoS with Neighbor Discovery,
- Can also throttle IPv6 traffic when processed in SW while IPv4 is in HW (legacy platform)
- In doubts: `show proc cpu | include IPv6`

```
class-map match-all ipv6
  match protocol ipv6

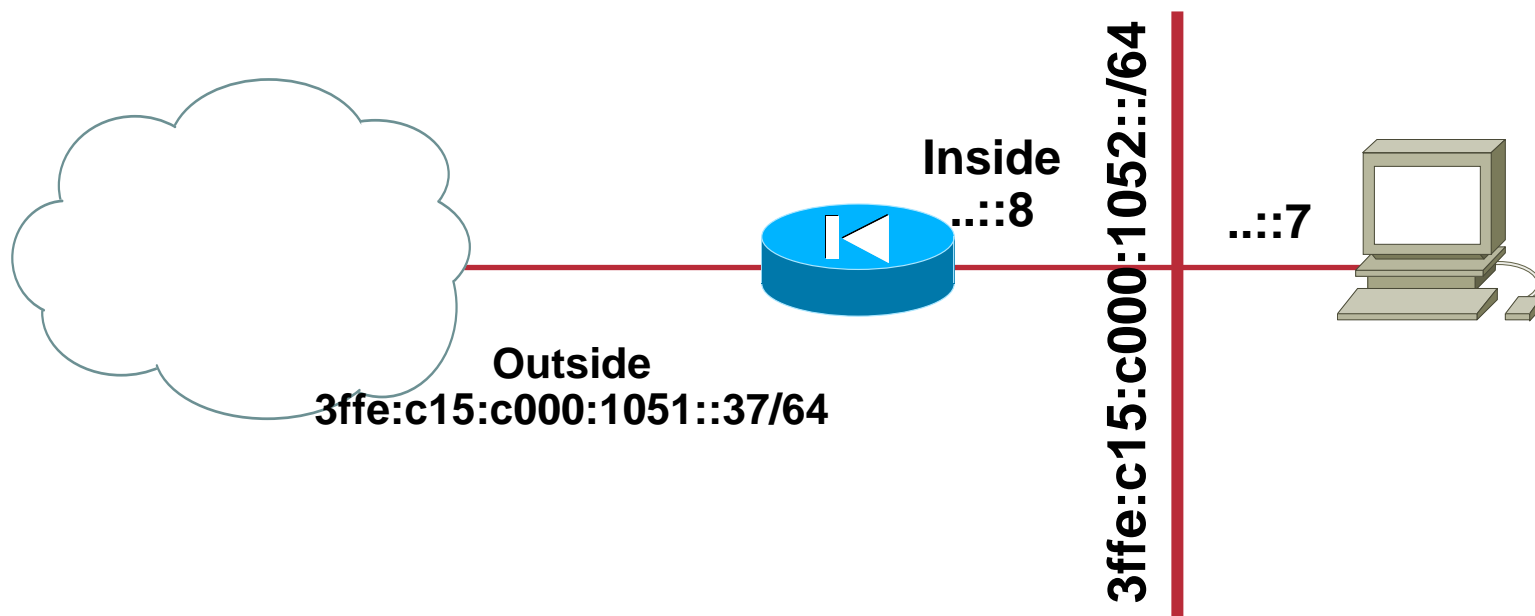
policy-map CoPP
  class ipv6
    police rate 100 pps
      conform-action transmit
      exceed-action drop

control-plane
  service-policy input CoPP
```

ASA and PIX Firewall IPv6 Support

- **Recognition of IPv6 traffic**
Dual-stack, IPv6 only, IPv4 only
- **Extended IP ACL with stateful inspection**
- **Application awareness**
HTTP, FTP, telnet, SMTP, TCP, SSH, UDP
- **uRPF**
- **v6 Frag guard**
- **IPv6 header security checks**
- **Management access via IPv6**
Telnet, SSH, HTTPS

ASA: Sample IPv6 Topology



ASA and PIX 7.0: ACL Very Similar to Cisco IOS

For Your Reference

```
interface Ethernet0
  nameif outside
  ipv6 address 3ffe:c15:c000:1051::37/64
  ipv6 enable
interface Ethernet1
  nameif inside
  ipv6 address 3ffe:c15:c000:1052::1/64
  ipv6 enable

ipv6 unicast-routing

ipv6 route outside ::/0 3ffe:c15:c000:1051::1

ipv6 access-list SECURE permit tcp any host
3ffe:c15:c000:1052::7 eq telnet
ipv6 access-list SECURE permit icmp6 any
3ffe:c15:c000:1052::/64

access-group SECURE in interface outside
```

ASA and PIX 7.0: Stateful Inspection

```
pixA# show conn
4 in use, 7 most used
ICMP out fe80::206:d7ff:fe80:2340:0 in
fe80::209:43ff:fea4:dd07:0 idle 0:00:00 bytes 16
UDP out 3ffe:c15:c000:1051::138:53 in
3ffe:c15:c000:1052::7:50118 idle 0:00:02 flags -
TCP out 2001:200:0:8002:203:47ff:fea5:3085:80 in
3ffe:c15:c000:1052::7:11009 idle 0:00:14 bytes 8975 flags
UfFRIO
TCP out 3ffe:c15:c000:1051::1:11008 in
3ffe:c15:c000:1052::7:23 idle 0:00:04 bytes 411 flags UIOB
```


“There is no reason anymore to let your site wide open for IPv6”

**An IPv6 Site Admin
Previously Fully Opened In IPv6
and Restricted in IPv4**

Enterprise Deployment: Secure IPv6



Secure IPv6 Traffic over IPv6 Public Network

- Since 12.4(6)T, IPsec also works for IPv6
- Using the Virtual Interface

```
interface Tunnel0
  no ip address
  ipv6 address 2001:DB8::2811/64
  ipv6 enable
  tunnel source Serial0/0/1
  tunnel destination 2001:DB8:7::2
  tunnel mode ipsec ipv6
  tunnel protection ipsec profile ipv6
```

Secure IPv6 over IPv4 Public Internet

- **How can we transport IPv6 securely over IPv4 Internet?**

No traffic sniffing

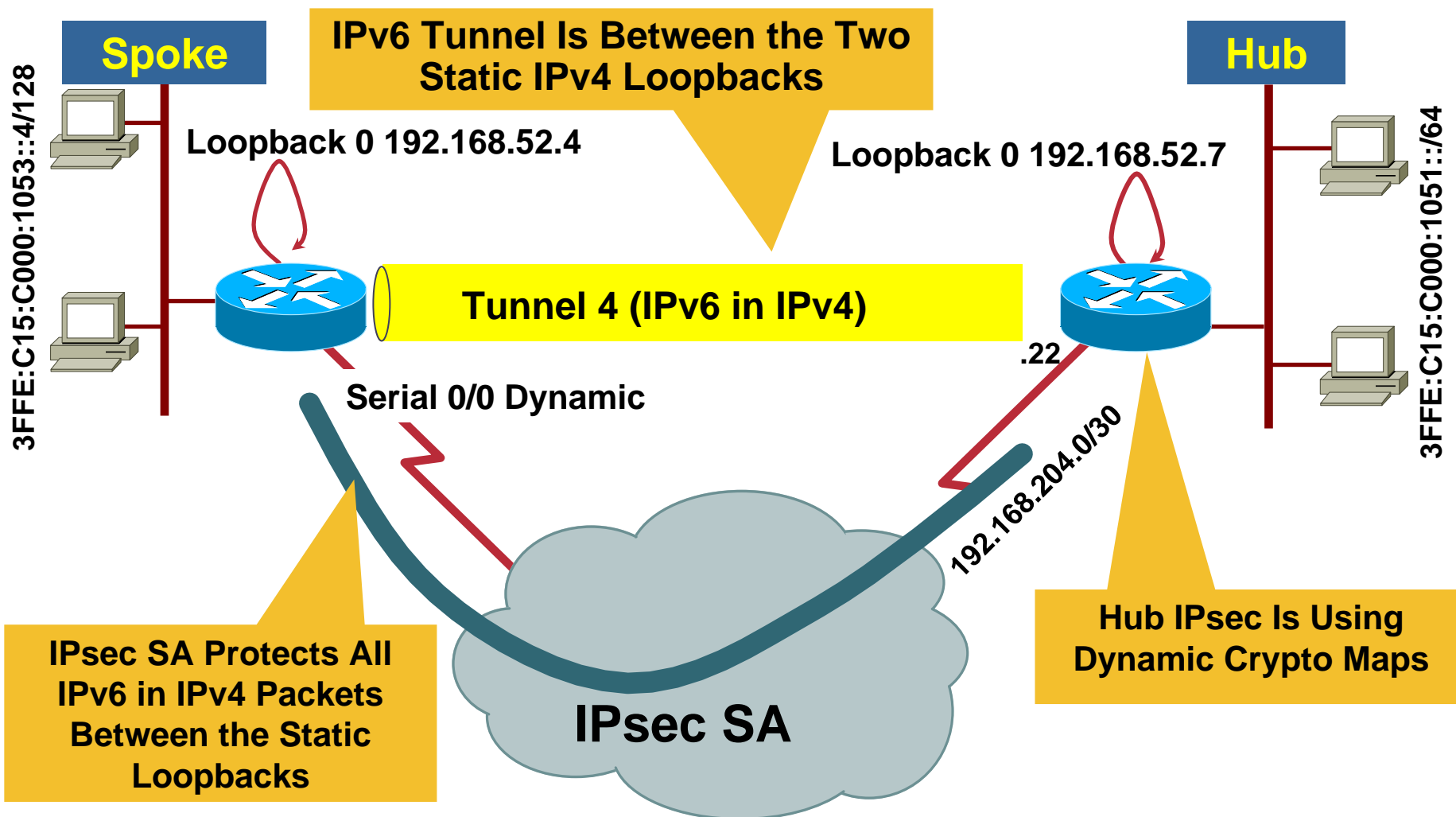
No traffic injection

- **IPsec**

Site to site: encrypting IPv6 tunnels

Remote access: encrypting ISATAP or IPv6 tunnels

Secure Site to Site IPv6 Connectivity



Key Design Points

- Requires a fixed IPv4 address for hub
- IPv6-in-IPv4 tunnels are **anchored** on IPv4 loopbacks

Tunnels requires **static** sources and destinations

- IPsec dynamic crypto maps are used

Allows for **dynamic** spoke IPv4 addresses

IPsec works on IPv4 packets (containing the IPv4 packets)

- Traffic initiated from spokes (hub is using dynamic crypto maps)

IPv6 for Remote Devices Solutions

- **Enabling IPv6 traffic inside the Cisco VPN Client tunnel**

NAT and Firewall traversal support

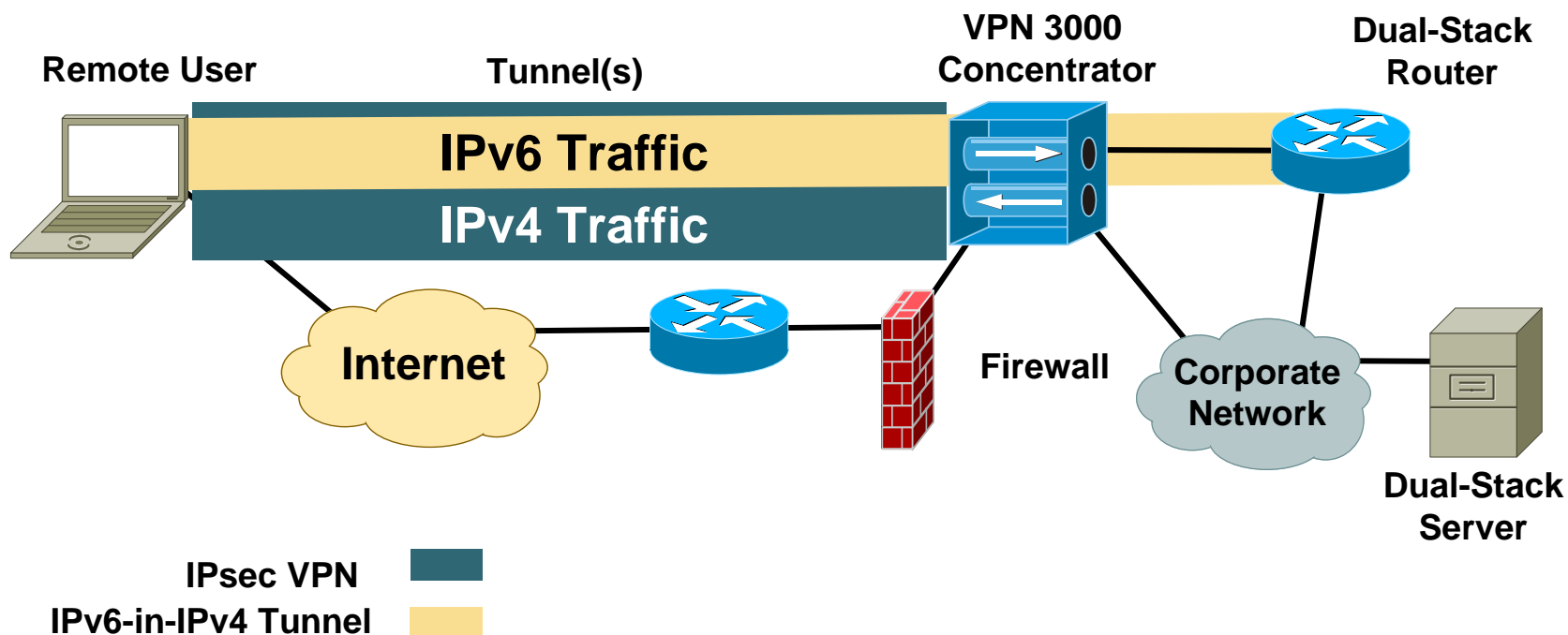
- **Allow remote host to establish a v6-in-v4 tunnel either automatically or manually**

ISATAP—Intra Site Automatic Tunnel Addressing Protocol

Configured—Static configuration for each side of tunnel

Fixed IPv6 address enables server's side of any application to be configured on an IPv6 host that could roam over the world

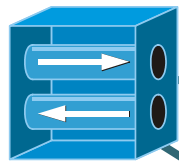
IPv6-in-IPv6 Tunnel Example



Note: The VPN concentrator could be replaced with a VPN-enabled Cisco IOS Router or PIX™

Router Configuration: ISATAP

VPN 3000 Concentrator

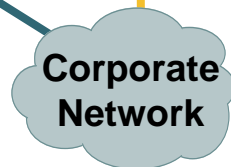


F0/0



Dual-Stack
Router

F0/1



Corporate
Network

Dual-stack router configuration

```

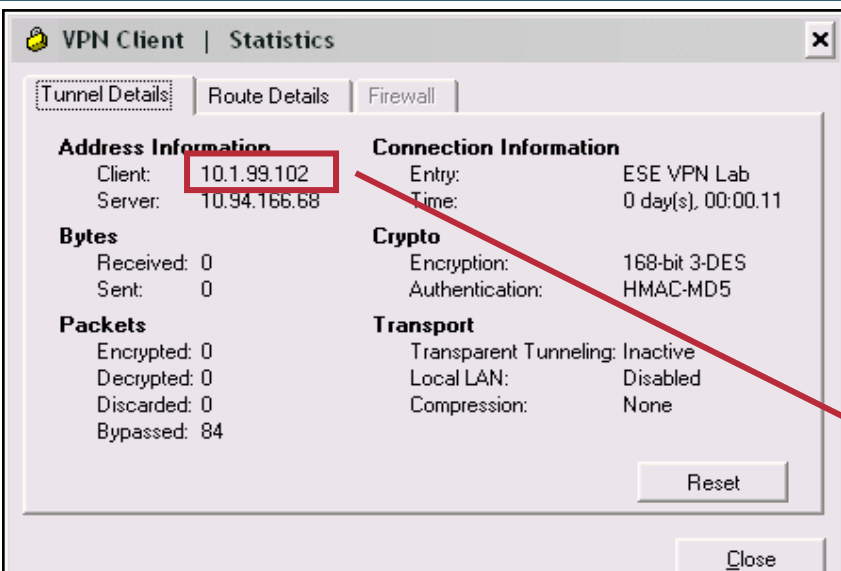
ipv6 unicast-routing
!
interface FastEthernet0/0
  description TO VPN 3000
  ip address 20.1.1.1 255.255.255.0
!
interface FastEthernet0/1
  description TO Campus Network
  ipv6 address 3FFE:C15:C003:111C::2/64
!
interface Tunnel0
  no ip address
  ipv6 address 3FFE:C15:C003:1101::/64
  eui-64
  no ipv6 nd suppress-ra
  tunnel source FastEthernet0/0
  tunnel mode ipv6ip isatap
  
```

ISATAP Address Format:

64-bit Unicast Prefix	32-bit	32-bit
0000:5EFE:	Interface ID	IPv4 Addr.

3ffe:c15:c003:1101:0:5efe:20.1.1.1

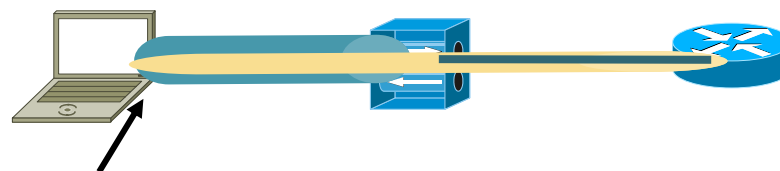
Does It Work?



Windows XP
Client

VPN 3000

Dual-Stack
Router



10.1.99.102 - VPN address

3ffe:c15:c003:1101:0:5efe:10.1.99.102 - IPv6 address

Interface 2: Automatic Tunneling Pseudo-Interface

Addr Type	DAD State	Valid Life	Pref. Life	Address
Public	Preferred	29d23h56m5s	6d23h56m5s	3ffe:c15:c003:1101:0:5efe:10.1.99.102
Link	Preferred	infinite	infinite	fe80::5efe:10.1.99.102

```
netsh interface ipv6>show route
Querying active state...
```

Publish	Type	Met	Prefix	Idx	Gateway/Interface Name
no	Autoconf	9	3ffe:c15:c003:1101::/64	2	Automatic Tunneling Pseudo-Interface
no	Manual	1	::/0	2	fe80::5efe:20.1.1.1

Conclusion



Summary Findings

- **IPv6 makes some things better, other things worse, and most things are just different, but no more or less secure**

Better

Automated scanning and worm propagation is harder due to huge subnets

Worse

Increased complexity in addressing and configuration

Lack of familiarity with IPv6 among operators

Vulnerabilities in transition techniques

- **Most of the legacy issues with IPv4 security remain in IPv6**

For example, ARP security issues in IPv4 are simply replaced with ND security issues in IPv6

Key Take Away

- **So, nothing really new in IPv6**
- **Security enforcement is possible**
Control your IPv6 traffic as you do for IPv4
- **Leverage IPsec to secure IPv6 when possible**

Q and A



Recommended Reading

- Continue your Networkers learning experience with further reading for this session from Cisco Press™
- Check the Recommended Reading flyer for suggested books

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Tuesday, June 20 at 12:15 p.m.

Wednesday, June 21 at 12:15 p.m.

Thursday, June 22 at 12:15 p.m. and 2:00 p.m.

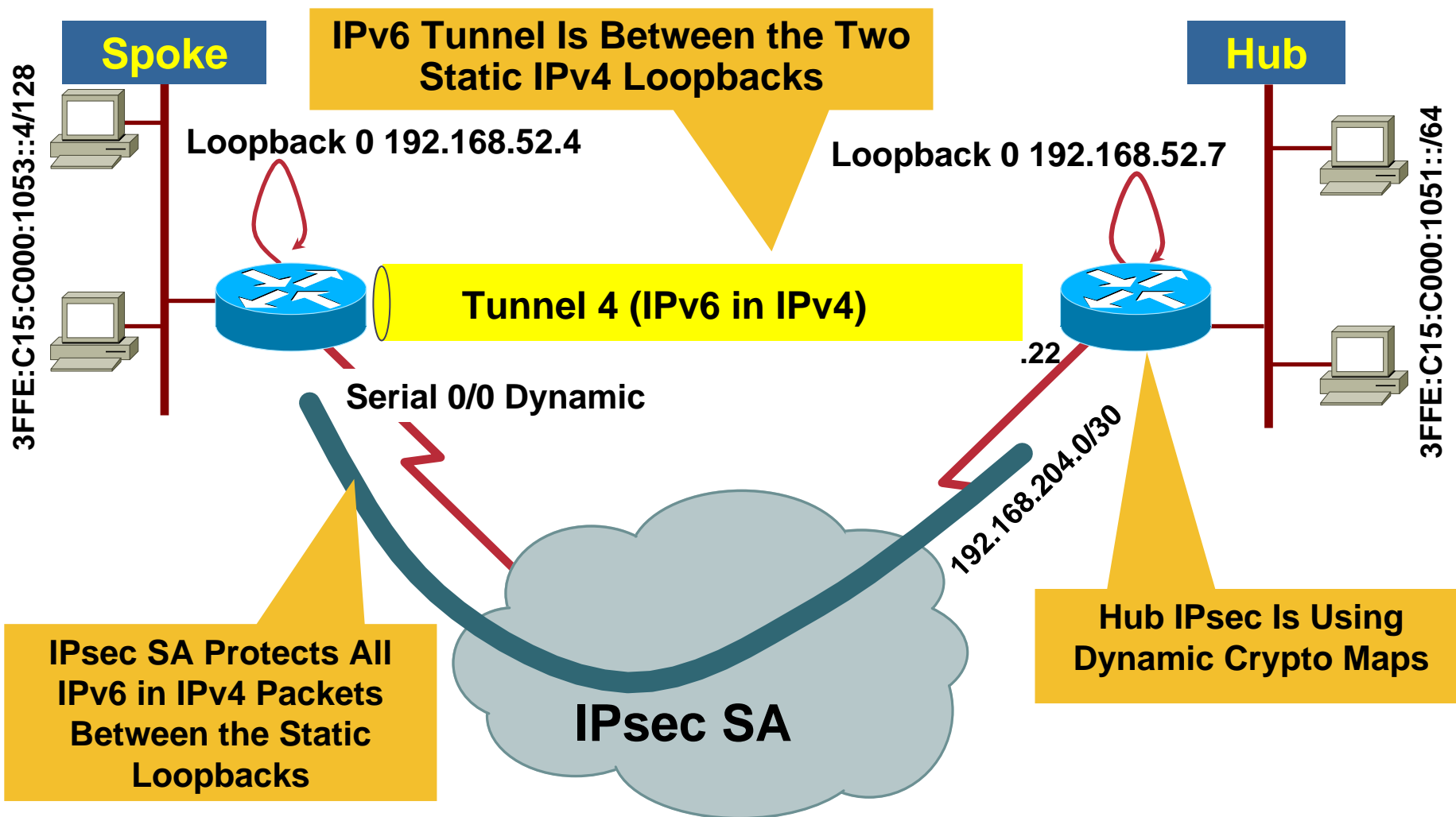




Reference Slides



Secure Site to Site IPv6 Connectivity



Spoke Configuration/1: IPv6 Tunnels

```
interface Loopback0
  ip address 192.168.52.4 255.255.255.255

interface Tunnel4
  no ip address
  ipv6 unnumbered FastEthernet0/0
  ipv6 enable
  tunnel source Loopback0
  tunnel destination 192.168.52.7
  tunnel mode ipv6ip
!
ip route 192.168.52.0 255.255.255.0 Serial0/0
```

**Static IPv4
Addresses**

Spoke Configuration/2: IPv4 IPsec

```
crypto ipsec transform-set 3DES esp-3des
!
```

```
crypto map IPV6_SEC 10 ipsec-isakmp
  set peer 192.168.204.26
  set transform-set 3DES
  match address SELECTOR
```

IPv4 Address of Hub

```
!
interface Serial10/0
  crypto map IPV6_SEC
```

IPsec Traffic Selectors:
Fixed IPv4 Loopback
Addresses, i.e.,
Encapsulated IPv6 Traffic

```
!
ip access-list extended SELECTOR
  permit 41 host 192.168.52.4 host 192.168.52.7
```

Hub Configuration/1: IPv6 Tunnels

```
interface Loopback0
  ip address 192.168.52.7 255.255.255.255
!
interface Tunnel4
  no ip address
  ipv6 unnumbered FastEthernet0/1
  ipv6 enable
  tunnel source Loopback0
  tunnel destination 192.168.52.4
  tunnel mode ipv6ip
```

**Static IPv4
addresses**

... a lot more interfaces Tunnel...

```
ip route 192.168.52.0 255.255.255.0 Serial0/0
```

Hub Configuration/2: IPv4 IPsec

```
crypto ipsec transform-set 3DES esp-3DES
!
crypto dynamic-map TEMPLATE 10
  set transform-set 3DES
  match address SELECTOR
!
crypto map IPV6_SEC 10 ipsec-isakmp dynamic TEMPLATE
!
interface Serial10/0
  ip address 192.168.204.26 255.255.255.255
  crypto map IPV6_SEC
!
ip access-list extended SELECTOR
  permit 41 host 192.168.52.7 192.168.52.0 0.0.0.255
```

**Dynamic crypto map:
Allow IPsec from every
IP address with correct
IKE authentication**

**IPsec traffic selectors:
fixed IPv4 loopback
addresses, i.e.,
encapsulated IPv6 traffic**